

# On maximum volume submatrices and cross approximation

Alice Cortinovis   Daniel Kressner   Stefano Massei

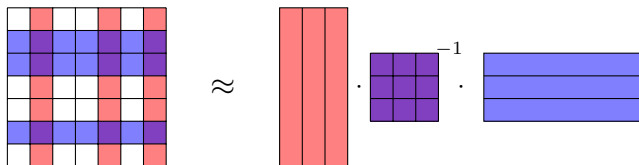
ANCHP  
Institute of Mathematics  
École Polytechnique Fédérale de Lausanne

Low-Rank Optimization and Applications  
Leipzig, 04 April 2019



# Low-rank approximation from rows and columns of $A$

Idea [Goreinov/Tyrtysnikov/Zamarashkin'1997, Bebendorf'2000, many others...]:  
Use selected rows and columns of  $A$  to build “cross approximation”



- Boundary element method [Bebendorf'2000]
- General tool for assembling  $\mathcal{H}$ -matrices,  $\mathcal{H}^2$ -matrices, ... [Hackbusch'2015]
- Uncertainty quantification [Harbrecht/Peters/ Schneider'2012]
- Kernel-based learning [Bach/Jordan'2005] and spectral clustering [Fowlkes/Belongie/Chung/Malik'2004] (Nyström method)
- Extension to low-rank tensor approximation [Oseledets/Tyrtysnikov'2010, Ballani/Grasedyck/Kluge'2013, Savostyanov'2014, ...]

- 1 Maxvol submatrices
  - Connection to low-rank approximation problem
  - NP hardness
- 2 Maxvol submatrices for structured matrices
  - symmetric positive semidefinite matrices;
  - diagonally dominant matrices.
- 3 Error bounds for cross approximation
  - general matrices;
  - symmetric positive semidefinite matrices;
  - (doubly) diagonally dominant matrices.
- 4 Cross approximation for functions

# Maxvol submatrices

# Quasi-optimal choice of rows and columns

How close is error of cross approximation

$$\text{error} = \begin{array}{|c|c|c|c|c|} \hline \color{red}{\square} & \color{red}{\square} & \color{red}{\square} & \color{red}{\square} & \color{red}{\square} \\ \hline \color{blue}{\square} & \color{blue}{\square} & \color{blue}{\square} & \color{blue}{\square} & \color{blue}{\square} \\ \hline \color{red}{\square} & \color{red}{\square} & \color{red}{\square} & \color{red}{\square} & \color{red}{\square} \\ \hline \color{blue}{\square} & \color{blue}{\square} & \color{blue}{\square} & \color{blue}{\square} & \color{blue}{\square} \\ \hline \color{red}{\square} & \color{red}{\square} & \color{red}{\square} & \color{red}{\square} & \color{red}{\square} \\ \hline \end{array} - \begin{array}{|c|c|c|} \hline \color{red}{\square} & \color{red}{\square} & \color{red}{\square} \\ \hline \color{red}{\square} & \color{red}{\square} & \color{red}{\square} \\ \hline \color{red}{\square} & \color{red}{\square} & \color{red}{\square} \\ \hline \end{array} \cdot \begin{array}{|c|c|} \hline \color{blue}{\square} & \color{blue}{\square} \\ \hline \color{blue}{\square} & \color{blue}{\square} \\ \hline \end{array}^{-1} \cdot \begin{array}{|c|c|c|} \hline \color{blue}{\square} & \color{blue}{\square} & \color{blue}{\square} \\ \hline \color{blue}{\square} & \color{blue}{\square} & \color{blue}{\square} \\ \hline \color{blue}{\square} & \color{blue}{\square} & \color{blue}{\square} \\ \hline \end{array}$$

to best rank- $m$  approximation error  $\min_{B,C \in \mathbb{R}^{n \times m}} \|A - BC^T\|_2 = \sigma_{m+1}(A)$ ?

Theorem ([Goreinov/Tyrtshnikov'2001])

If  $\begin{array}{|c|c|c|} \hline \color{blue}{\square} & \color{blue}{\square} & \color{blue}{\square} \\ \hline \color{blue}{\square} & \color{blue}{\square} & \color{blue}{\square} \\ \hline \color{blue}{\square} & \color{blue}{\square} & \color{blue}{\square} \\ \hline \end{array}$  is  $m \times m$  submatrix of maximum volume (maximum absolute value of the determinant), then

$$\|\text{error}\|_{\max} \leq (m+1)\sigma_{m+1}(A),$$

where  $\|B\|_{\max} = \max_{i,j=1,\dots,n} |b_{ij}|$ .

*Proof idea by [Goreinov/Tyrtshnikov'2001]. Assume w.l.o.g.  $n = m + 1$ . Then*

$$\|A^{-1}\|_2 = \frac{1}{\sigma_{m+1}(A)}.$$

On the other hand,

$$A^{-1} = \frac{1}{\det A} \text{adj}(A)$$

implies that maximum element of  $A^{-1}$  is at entry  $(m + 1, m + 1)$ . In turn,

$$\|\text{error}\|_{\max} = |(A^{-1})_{m+1,m+1}| = \|A^{-1}\|_{\max}.$$

Hence,

$$\sigma_{m+1}(A)^{-1} = \|A^{-1}\|_2 \leq (m + 1) \|A^{-1}\|_{\max} = (m + 1) |(A^{-1})_{m+1,m+1}|.$$

Remark.

- Maxvol approximation is actually optimal in some weird sense. Consider mixed norm

$$\|B\|_{\infty \rightarrow 1} := \sup \|Bx\|_1 / \|x\|_{\infty}, \quad \|B\|_{1 \rightarrow \infty} := \sup \|Bx\|_{\infty} / \|x\|_1 = \|B\|_{\max}.$$

and approximation numbers

$$\beta_{k+1}(A) := \min\{\|E\|_{\infty \rightarrow 1} : \text{rank}(A + E) \leq k\}, \quad k = 0, \dots, n-1.$$

Then  $\beta_n(A) = \|A^{-1}\|_{\max}^{-1}$  see [Higham'2002]. An extension of the proof shows

$$\|\text{error}\|_{\max} = \beta_{m+1}(A).$$

By norm equivalence,

$$\|\text{error}\|_{\max} \leq (m+1)^2 \min\{\|E\|_{\max} : \text{rank}(A + E) \leq m\}.$$

and the constant is tight. Recovers result by [Goreinov/Tyrtysnikov'2011].

## (Approximate) maxvol is NP hard

- Papadimitriou'1984 reduced NP-complete SAT via a subgraph selection problem to maxvol for 0/1 matrices  $\rightsquigarrow$  maxvol is NP hard.
- Di Summa et al.'2014 showed that it is NP-hard to approximate maxvol within factor that does not grow at least exponentially with  $m$ .

*Remark.* Due to nature of determinants, second result has little implication for low-rank approximation. Consider

$$A = \begin{bmatrix} I_m & 0 \\ 0 & B_m \end{bmatrix}, \quad B_m = \text{tridiag}\left[\frac{1}{2}, 1, -\frac{1}{2}\right]$$

Note that

$$|\det(I_m)| = 1, \quad |\det(B_m)| \sim \left(\frac{1 + \sqrt{2}}{2}\right)^m.$$

Nevertheless, the first  $m$  rows/columns constitute an excellent choice for cross approximation:

$$\|\text{error}\|_{\max} = \|B_m\|_{\max} = 1 = \sigma_{m+1}(B).$$

Is it NP hard to find cross approximation with polynomial constant?



# Maxvol submatrices for structured matrices

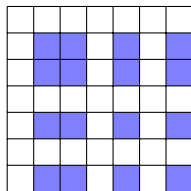
# Maxvol submatrix for SPSD

Let  $A$  be symmetric positive semidefinite (SPSD).

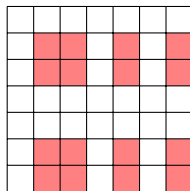
Obvious: Element of maximum absolute value of  $A$  on the diagonal.

Less obvious: Submatrix of maximum volume of  $A$  can always be chosen to be **principal**.

Principal submatrix



Not principal




# Maxvol submatrix for SPSD

Theorem ([Cortinovis/K./Massei'2019])



If  $A$  is symmetric positive semidefinite then the maximum volume  $m \times m$  submatrix is attained by a principal submatrix.

Volume of a (rectangular) matrix = product of its singular values.

① Cholesky decomposition  $A = C^T \cdot C$ : 

② Inequality for product of singular values of a product [Horn/Johnson'1991]:

$$\text{Volume of } \text{purple square} \leq \text{Volume of } \text{blue rectangle} \cdot \text{Volume of } \text{red rectangle}$$

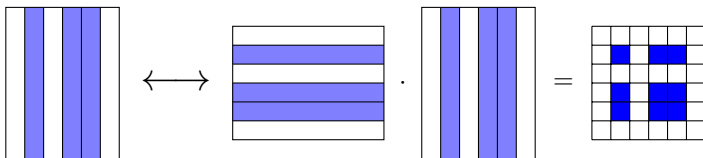
③ Relate volume of  and  to volume of principal submatrices and conclude. □

# Consequences

- 1 One-to-one correspondence with column selection problem:

*Given  $n \times k$  matrix  $B$ , find  $n \times m$  submatrix of maximum volume.*

Also arises in low-rank approximation problems [Deshpande/Rademacher/Vempala/Wang'2006] and rank-revealing LU factorizations [Pan'2000].



- 2 Column selection problem is NP-hard [Çivril/Magdon-Ismail'2009]  $\rightsquigarrow$  maximum volume submatrix problem is NP-hard even when restricted to symmetric positive semidefinite matrices.

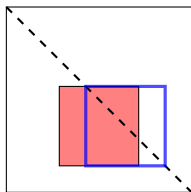
# Maxvol submatrix for diagonally dominant

## Definition

$A \in \mathbb{R}^{n \times n}$  is (row) diagonally dominant if  $\sum_{j=1, \dots, n; j \neq i} |a_{ij}| \leq |a_{ii}|$  for all  $i = 1, \dots, n$ .

## Theorem ([Cortinovis/K./Massei'2019])

*If  $A$  is diagonally dominant then the maxvol  $m \times m$  submatrix is attained by a principal submatrix.*



*Proof idea.*

- 1 Reduction to upper triangular case via LU factorization.

$$\begin{aligned} A(I, J) &= \begin{array}{|c|} \hline \text{red square} \\ \hline \end{array} = \begin{array}{|c|} \hline \text{purple triangle} \quad 0 \\ \hline \end{array} \cdot \begin{array}{|c|} \hline \text{red square} \\ \hline 0 \end{array} \\ A(I, I) &= \begin{array}{|c|} \hline \text{blue square} \\ \hline \end{array} = \begin{array}{|c|} \hline \text{purple triangle} \quad 0 \\ \hline \end{array} \cdot \begin{array}{|c|} \hline \text{blue triangle} \\ \hline 0 \end{array} \end{aligned}$$

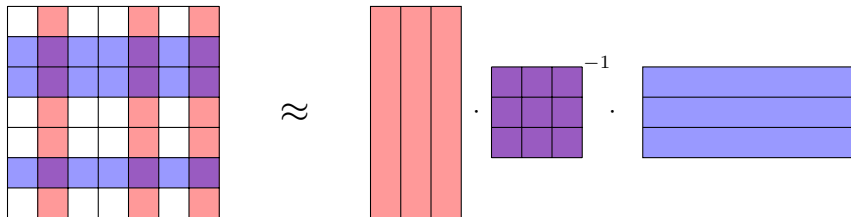
- 2 Factor  $U$  inherits diagonal dominance from  $A$ ! For unit upper triangular diagonally dominant matrix: Each submatrix has  $|\det| \leq 1$ . □

Known special case: for  $m = n - 1$ , the result of the theorem is covered in the proof of Theorem 2.5.12 in [Horn/Johnson'1991].

# Cross approximation

# Cross approximation algorithm with full pivoting

Aim: Finding quasi-optimal row/column indices.



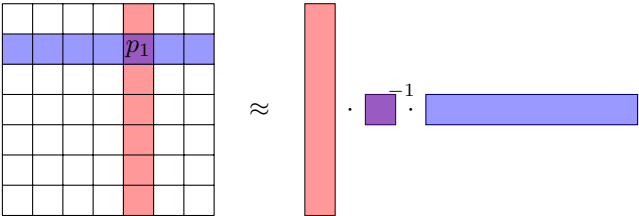


# Cross approximation algorithm with full pivoting

ACA = Adaptive Cross Approximation [Bebendorf'2000, Carvajal/Chapman/ Geddes'2005, ...]

Greedy algorithm for volume maximization.

**First step:** Select element of maximum absolute value,  $p_1$  (first pivot).



Denote

$$R_1 := \begin{matrix} \square & \square & \square & \square & \square & \square \\ \color{blue}\square & \color{blue}\square & \color{blue}\square & p_1 & \color{blue}\square & \color{blue}\square \\ \square & \square & \square & \square & \square & \square \\ \square & \square & \square & \square & \square & \square \\ \square & \square & \square & \square & \square & \square \\ \square & \square & \square & \square & \square & \square \end{matrix} - \begin{matrix} \square \\ \square \\ \square \\ \square \\ \square \\ \square \end{matrix} \cdot \begin{matrix} \color{purple}\square^{-1} \\ \cdot \\ \cdot \end{matrix} \cdot \begin{matrix} \square & \square & \square & \square & \square & \square \\ \square & \square & \square & \square & \square & \square \\ \square & \square & \square & \square & \square & \square \\ \square & \square & \square & \square & \square & \square \\ \square & \square & \square & \square & \square & \square \\ \square & \square & \square & \square & \square & \square \end{matrix} = \begin{matrix} \square & \square & \square & \square & \square & \square \\ 0 & 0 & 0 & 0 & 0 & 0 \\ \square & \square & \square & \square & \square & \square \\ \square & \square & \square & \square & \square & \square \\ \square & \square & \square & \square & \square & \square \\ \square & \square & \square & \square & \square & \square \end{matrix}$$

# Cross approximation algorithm with full pivoting

Situation after  $k - 1$  steps:

	0			0		
0	0	0	0	0	0	0
	0			0		
	0			0		
	0			0		
0	0	0	0	0	0	0
	0			0		

 $= R_{k-1} =$ 


 $-$ 


 $\cdot$ 


 $^{-1}$  $\cdot$ 


# Cross approximation algorithm with full pivoting

Situation after  $k - 1$  steps:

The diagram shows the matrix  $R_{k-1}$  as a product of three matrices:

- A permutation matrix (left): A 6x6 grid with red vertical bars in columns 1, 2, 5, and 6, and blue horizontal bars in rows 3 and 6. The element at row 3, column 5 is labeled  $p_3$ .
- An equals sign.
- The matrix  $R_{k-1}$  (middle): A 6x6 grid with red vertical bars in columns 1 and 4, and blue horizontal bars in rows 2 and 5. The element at row 2, column 4 is labeled  $p_1$ , and the element at row 5, column 1 is labeled  $p_2$ .
- An equals sign.
- A permutation matrix (right): A 6x6 grid with red vertical bars in columns 1 and 2.
- A dot.
- A 2x2 block matrix (right): A 2x2 grid with purple cells. The top-right cell is labeled  $-1$ .
- A dot.
- A matrix (right): A 6x6 grid with blue horizontal bars in rows 2 and 5.

**At  $k$ -th step:** Choose element of maximum absolute value in remainder  $R_{k-1}$  ( $k$ -th pivot  $p_k$ ).

# Cross approximation algorithm with full pivoting

Situation after  $k - 1$  steps:

	0			0			
0	0	0	0	0	0	0	0
0				0			$p_3$
0				0			
0				0			
0	0	0	0	0	0	0	0
0				0			

 $= R_{k-1} =$ 

							$p_1$
							$p_2$

 $-$ 


 $\cdot$ 


 $^{-1}$ 


At  $k$ -th step: Choose element of maximum absolute value in remainder  $R_{k-1}$  ( $k$ -th pivot  $p_k$ ).

	0			0			
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0
	0			0			
	0			0			
0	0	0	0	0	0	0	0
	0			0			

 $= R_k =$ 

							$p_1$
							$p_3$
							$p_2$

 $-$ 


 $\cdot$ 


 $^{-1}$ 


# Equivalence to Gaussian elimination and LDU factorization

- Cross approximation = greedy algorithm for volume maximization.
- Equivalent to Gaussian elimination with complete pivoting [Bebendorf'2000]. In particular, no need to compute inverses at each step!
- (Up to permutations) we obtain an incomplete LDU factorization

$$A - R_m = L_m D_m U_m = \begin{array}{|c|} \hline \text{0} \\ \hline \text{red triangle} \\ \hline \end{array} \cdot \begin{array}{|c|} \hline \text{0} \\ \hline \text{purple diagonal} \\ \hline \text{0} \\ \hline \end{array} \cdot \begin{array}{|c|} \hline \text{blue triangle} \\ \hline \text{0} \\ \hline \end{array}$$

where

- $D_m = \text{diag}(p_1, \dots, p_m)$  contains pivot elements;
- $L_m$  and  $U_m$  have ones on the diagonal.

## Error of cross approximation

Analysis of error  $\|R_m\|_{\max}$  obtained after  $m$  steps of cross approximation: By performing one additional step

$$\|R_m\|_{\max} = \|\text{Schur complement of } A_{11} \text{ in } A\|_{\max} = |p_{m+1}|.$$

Ideally  $|p_{m+1}|$  is close to  $\sigma_{m+1}(A)$ .

Wlog, restrict to  $(m+1) \times (m+1)$  matrices and consider factorization  $A = LDU$ :

$$\frac{1}{\sigma_{m+1}(A)} = \|A^{-1}\|_2 \leq \|U^{-1}\|_2 \|D^{-1}\|_2 \|L^{-1}\|_2.$$

What can go wrong?

- 1 Intermediate pivots can be  $\ll |p_{m+1}| \rightsquigarrow \|D^{-1}\|_2 \not\ll \frac{1}{|p_{m+1}|}$ .
- 2  $\|L^{-1}\|_2$  and  $\|U^{-1}\|_2$  can be large.

Closely related to but **not** covered by numerical linear algebra literature on error analysis of (rank-revealing) LU decompositions.

## Error bounds for the general case

- 1 Consider growth factor =  $\rho_k := \sup \left\{ \frac{\|R_k\|_{\max}}{\|A\|_{\max}} \mid \text{rank}(A) \geq k \right\}$ , playing prominent role in error analysis of LU factorization. [Wilkinson'1961] proved for complete pivoting

$$\rho_k \leq 2\sqrt{k+1}(k+1)^{\ln(k+1)/4}.$$

Note that bound is not tight; usually  $\rho_k = O(1)$ . Obtain

$$\|D^{-1}\|_2 \leq \frac{\rho_m}{|p_{m+1}|}.$$

- 2  $L$  and  $U$  have ones on the diagonal and all other entries have absolute value  $\leq 1$  because of full pivoting  $\rightsquigarrow \|L^{-1}\|_2 \leq 2^m$  and  $\|U^{-1}\|_2 \leq 2^m$ ; see [Higham'1987].

Theorem ([Cortinovis/K./Massei'2019])

After  $m$  steps, error of cross approximation satisfies

$$\|R_m\|_{\max} \leq 4^m \rho_m \cdot \sigma_{m+1}(A).$$

# Symmetric positive semidefinite case

Benefits for SPSD matrices:

- ① **Huge:** Diagonal pivoting is sufficient (SPSD preserved by Schur compl)
- ② Minor: Pivots do not grow  $\rightsquigarrow \rho_m$  replaced by 1 in the theorem.

## Corollary

If  $A$  is symmetric positive semidefinite then

$$\|R_m\|_{\max} \leq 4^m \cdot \sigma_{m+1}(A).$$

This matches a result of [Harbrecht/Peters/Schneider'2012].

Bound is tight for SPSD. Kahan'1966:

$$U = L^T = \begin{bmatrix} 1 & -1 & \cdots & -1 \\ 0 & 1 & \ddots & \vdots \\ \vdots & \ddots & \ddots & -1 \\ 0 & \cdots & 0 & 1 \end{bmatrix} \in \mathbb{R}^{(m+1) \times (m+1)}$$

$$\|A^{-1}\|_2 = \|U^{-1}\|_2^2 \sim 4^m \rightsquigarrow \sigma_{m+1}(A) \sim 4^{-m}$$

On the other hand,  $|p_{m+1}| = 1$ .



# Diagonally dominant case

Benefits for diagonally dominant matrices:

- 1 Diagonal pivoting is sufficient (diagonal dominance preserved by Schur compl)
- 2 Small growth factor:  $\|R_k\|_{\max} \leq 2\|A\|_{\max}$  for every  $k$ . [Wilkinson'1961]
- 3 In the LDU factorization,  $U$  is diagonally dominant. Hence,  $\|U^{-1}\|_2 \leq m$ . [Peña'2004]

## Corollary

If  $A$  is diagonally dominant then  $\|R_m\|_{\max} \leq (m+1) \cdot 2^{m+1} \cdot \sigma_{m+1}(A)$ .

- 4 If  $A$  is doubly diagonally dominant (that is,  $A$  and  $A^T$  are diag. dom.) then also  $L^T$  is diagonally dominant.

## Corollary

If  $A$  is doubly diagonally dominant then  $\|R_m\|_{\max} \leq 2(m+1)^2 \cdot \sigma_{m+1}(A)$ .

Result relevant in spectral clustering based on cross approximation.

## Diagonally dominant case: Tightness of bounds

- Diagonally dominant case: example with

$$\frac{\|R_m\|_{\max}}{\sigma_{m+1}(A)} = \Theta(m^2).$$

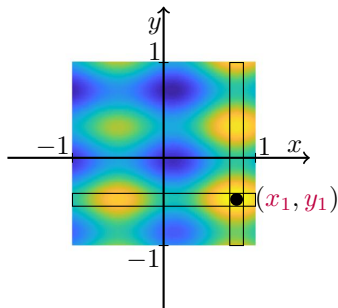
Related to studies on stability of LDU factorizations [e.g. Demmel/Koev'2004, Dopico/Koev'2011, Barreras/Peña'2012/13].

- Doubly diagonally dominant case: example with

$$\frac{\|R_m\|_{\max}}{\sigma_{m+1}(A)} = \Theta(m).$$

## Extension to functions

Consider bivariate function  $f : [-1, 1] \times [-1, 1] \rightarrow \mathbb{R}$  and choose point  $(x_1, y_1)$  of maximum absolute value.



“Rank-1 approximation” of  $f$  is separable function

$$f_1(x, y) = f(x, y_1) \cdot \frac{1}{f(x_1, y_1)} \cdot f(x_1, y).$$

Next steps  $\rightsquigarrow$  analogous to matrix algorithm.

[Bebendorf'2000, Carvajal/Chapman/Geddes'2005, Townsend/Trefethen'2015]

# Error bounds for functional approximation

## Definition

Bernstein ellipse  $\mathcal{E}_r =$  ellipse with foci  $\pm 1$  and sum of semiaxes  $r$ .

## Theorem ([Cortinovis/K./Massei'2019])

Assume that  $f(\cdot, y)$  admits an analytic extension  $\tilde{f}$  to the Bernstein ellipse  $\mathcal{E}_{r_0}$  for each  $y \in [-1, 1]$ . Choose  $1 < r < r_0$ . Denote

$$M := \sup_{\eta \in \partial \mathcal{E}_r, \xi \in [-1, 1]} |\tilde{f}(\eta, \xi)|.$$

Then the error after  $m$  steps satisfies

$$\|\text{error}_m\|_{\max} \leq \frac{2M\rho_m}{1 - 1/r} \cdot \left(\frac{r}{4}\right)^{-m}.$$

**Idea of proof:** error bound for cross approximation for general matrices applied to matrix interpolating the function in suitable points + standard polynomial approximation arguments.

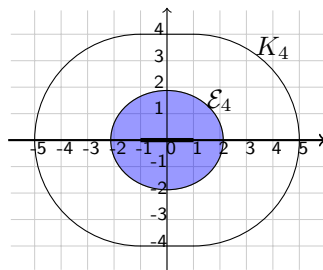
## Comparison to existing convergence results

$\rho_m$  has subexponential growth  $\rightsquigarrow$  Algorithm converges linearly with rate  $\frac{4}{r}$  for  $r > 4$ .

Previous convergence results for complete pivoting [Townsend/Trefethen'2015]: they need the functions  $f(\cdot, y)$  to have an analytic extension in

$$K_r := \{\text{points at distance } \leq r \text{ from the segment } [-1, 1]\}$$

for linear convergence with rate  $\frac{4}{r}$ .



# Conclusions

New results:

- 1 Maxvol submatrix of symmetric positive semidefinite or diagonally dominant matrices attained by principal submatrix.
- 2 Error analysis of cross approximation (both for matrix and function case).

Major open problem:

- 3 Cross approximation with complete pivoting always works well in practice. Find appropriate framework that captures this!

Next step:

- 4 Tensors!

Techreport available from [anchp.epfl.ch](http://anchp.epfl.ch):

- A. Cortinovis, DK, and S. Massei. On maximum volume submatrices and cross approximation for symmetric semidefinite and diagonally dominant matrices. February 2019.