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R-function Related to Entanglement of  
Formation

by

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# *R*-function Related to Entanglement of Formation

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By investigating the convex property of the function  $R$ , appeared in computing the entanglement of formation for isotropic states in Phys. Rev. Lett. **85**, 2625 (2000), and a tight lower bound of entanglement of formation for arbitrary bipartite mixed states in Phys. Rev. Lett. **95**, 210501 (2005), we show analytically that the very nice results in these papers are valid not only for dimensions 2 and 3 but any dimensions.

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Quantum entangled states are playing fundamental roles in quantum information processing [1]. The entanglement of formation (EOF) is a most important measure in quantifying the degree of entanglement [2, 3]. Considerable efforts have been spent on deriving EOF or its lower bound through analytical and numerical approaches ([4, 5] and references therein). In [4] the EOF for isotropic states are presented. It is shown that for  $F \geq 1/d$ , the EOF for isotropic states  $\rho_F$  is  $E(\rho_F) = co(R(F))$ , where “ $co(R)$ ” stands for the convex hull, the largest convex function that is bounded above by the function  $R$ . An explicit expression of  $co(R(F))$  has been derived for dimensions  $d = 2, 3$ , and its general form is conjectured for arbitrary  $d$ . While in [5] a tight lower bound of EOF for arbitrary bipartite mixed states is given, by using the property of  $R$  function derived in [4]. For any bipartite  $m \otimes n$  ( $m \leq n$ ) mixed quantum state  $\rho$ , the EOF  $E(\rho)$  satisfies  $E(\rho) \geq co(R(\Lambda))$ . In both cases the results depend on the convex property of the function  $R$ . These results are correct when the second derivative of  $R$  has one zero point, which has been shown to be true if the dimension is 2 or 3.

The  $R(\Lambda)$ -function ( $\Lambda$  corresponds to  $dF$  in  $R(F)$  [4]) has the form

$$R(\Lambda) = H_2(\gamma(\Lambda)) + (1 - \gamma(\Lambda)) \log_2(m - 1), \quad (1)$$

where

$$\gamma(\Lambda) = \frac{1}{m^2} \left( \sqrt{\Lambda} + \sqrt{(m-1)(m-\Lambda)} \right)^2 \quad (2)$$

and  $H_2(\cdot)$  is the standard binary entropy function,  $\Lambda \in [1, m]$ . For  $m \geq 4$ , it is conjectured that the second derivative of  $R$  w.r.t.  $\Lambda$  has still only one zero point, by numerically calculating the function  $R$  for a given  $m$ . In deed for  $m = 4$ , one can easily see this is true by plotting  $R(\Lambda)$ . In the following we show analytically that the results in both papers valid for arbitrary dimensions  $m \geq 5$ .

For simplicity we replace  $\log_2$  in (1) by the natural log. Without confusion we still use the notion  $R(\Lambda)$  below, which, in fact, differs a positive factor  $\log_2 e$  from the  $R(\Lambda)$  above. We first prove that there is one and only one point  $\Lambda_0$  between 1 and  $m-1$  such that  $R''(\Lambda_0) = 0$  for  $m \geq 5$ . Then we further show that there is no more zero points for  $R''(\Lambda)$  between  $m-1$  and  $m$ . By direct calculation we have the second derivative of  $R$  w.r.t  $\Lambda$

$$R''(\Lambda) = \gamma''(\Lambda) \log \frac{1 - \gamma(\Lambda)}{(m-1)\gamma(\Lambda)} - \frac{1}{\Lambda(m-\Lambda)}, \quad (3)$$

where

$$\gamma''(\Lambda) = -\frac{\sqrt{m-1}}{2} (\Lambda(m-\Lambda))^{-3/2}. \quad (4)$$

From which we get  $R''(1) = \lim_{\epsilon \rightarrow 0} R''(1+\epsilon) = +\infty$ . On the other hand,

$$R''(m-1) = -\frac{1}{m-1} \left( \log \frac{m-2}{2(m-1)} + 1 \right),$$

which is less than 0 for  $m \geq 5$ . Therefore for  $m \geq 5$  there exists  $\Lambda_0 \in (1, m-1)$  such that  $R''(\Lambda_0) = 0$ . From (3) and (4)  $\Lambda_0$  is the solution of  $g(\Lambda) = f(\Lambda)$ , where

$$g(\Lambda) = \log \frac{1 - \gamma(\Lambda)}{(m-1)\gamma(\Lambda)}, \quad f(\Lambda) = -2\sqrt{\frac{\Lambda(m-\Lambda)}{m-1}}.$$

As  $g'(\Lambda) > 0$ ,  $g(\Lambda)$  is a monotonically increasing function taking values from  $g(1) \rightarrow -\infty$  to

$$g(m-1) = 2 \log \frac{m-2}{2(m-1)} > -2.$$

While  $f(1) = f(m-1) = -2$ ,  $f''(\Lambda) > 0$ , i.e.  $f$  is convex. Therefore there is one and only one solution  $\Lambda_0$  to the equation  $g(\Lambda) = f(\Lambda)$  for  $\Lambda \in (1, m-1)$ .

We now show that there are no more solutions to  $R''(\Lambda) = 0$  for  $\Lambda \in (m-1, m)$ , i.e.  $R''(m-1+\delta) \neq 0$ ,  $\forall \delta \in (0, 1)$ . From (2), (3) and (4) this is equivalent to show  $F(\delta) \equiv \frac{1}{2}B(\delta) \log A(\delta) \neq -1$ , where,

$$B(\delta) = \sqrt{\frac{m-1}{(m-1+\delta)(1-\delta)}}, \quad A(\delta) = \frac{(mC(\delta))^2 - 1}{m-1},$$

and  $C(\delta) = (\sqrt{m-1+\delta} + \sqrt{(m-1)(1-\delta)})^{-1}$ . It is straight forward to verify that  $A(0) > 0$ . As the derivative of  $C(\delta)$  w.r.t  $\delta$ ,  $C'(\delta) > 0$ , we have  $A'(\delta) > 0$ . Hence  $\log A(\delta)$  increases as  $\delta$  increases. Similarly, as the derivative of  $(m-1)/((m-1+\delta)(1-\delta))$  w.r.t  $\delta$  is positive,  $B(\delta)$  also increases as  $\delta$  increases. Therefore  $F(\delta)$  is an increasing function of  $\delta$ . Moreover  $F(0) = \log(m-2)/(2(m-1)) \geq \log 3/8 > -1$ . We have  $F(\delta) \geq F(0) > -1, \forall \delta \in (0, 1)$  and  $m \geq 5$ . Thus  $R''(\Lambda) = 0$  has no solutions for  $\Lambda \in (m-1, m)$ .

We have shown that the  $R$ -function has only one reflection point for any dimensions. Therefore the construction of the largest convex function that is bounded above by the  $R$ -function in [4] is correct. Either the EOF for isotropic states in [4] and the tight lower bound of EOF in [5] are valid for arbitrary dimensions.

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