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by

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In this paper, we study entropic uncertainty relations on a finite-dimensional Hilbert space and provide several tighter bounds for multi-measurements, with some of them also valid for Rényi and Tsallis entropies besides the Shannon entropy. We employ majorization theory and actions of the symmetric group to obtain an *admixture bound* for entropic uncertainty relations for multi-measurements. Comparisons among all bounds for multi-measurements are given in two figures.

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I. INTRODUCTION

The most revolutionary departure of quantum mechanics from classical mechanics is that it is impossible to simultaneously measure two complementary variables of a particle in precision. Kennard's form of the Heisenberg uncertainty principle [1] displays vividly such an inequality for the standard deviation of position and momentum of a particle: $\sigma_Q \sigma_P \geq \frac{1}{2}$, where the Planck constant is taken as $\hbar = 1$. The corresponding entropic uncertainty of Białynicki-Birula and Mycielski [2] says that $h(Q) + h(P) \geq \log(e\pi)$, where Q and P stand for position and momentum respectively while h is the differential entropy: $h(Q) = -\int_{-\infty}^{\infty} f(x) \log f(x) dx$ with $f(x)$ being the probability density corresponding to Q .

In the seminal paper [3], Deutsch studied the entropic uncertainty relations on finite d -dimensional Hilbert spaces in terms of the Shannon entropy for any two measurements M_1 and M_2 (base 2 log is used unless stated otherwise):

$$H(M_1) + H(M_2) \geq -2 \log \frac{1 + \sqrt{c_1}}{2}, \quad (1)$$

where c_1 is the largest element in the overlap matrix $c(M_1, M_2)$ of the two measurements. Later Maassen and Uffink [4, 5] derived the influential generalized quantum mechanical uncertainty relation which amounts to a tighter lower bound than Eq. (1). Recently Coles and Piani [6] proved that, for any two measurements

$M_j = \{|u_{i_j}^j\rangle\}$ on a quantum state ρ over a finite dimensional Hilbert space

$$H(M_1) + H(M_2) \geq -\log c_1 + \frac{1 - \sqrt{c_1}}{2} \log \frac{c_1}{c_2}, \quad (2)$$

where c_2 is the second largest value among all overlaps $c(u_{i_1}^1, u_{i_2}^2) = |\langle u_{i_1}^1 | u_{i_2}^2 \rangle|^2$. Then Maassen-Uffink's bound is simply obtained by dropping the second term in RHS of Eq. (2).

More recently, S. Liu *et al.* [7] generalized Coles and Piani's method to give a lower bound for N measurements M_i :

$$\sum_{m=1}^N H(M_m) \geq -\log b + (N-1)S(\rho), \quad (3)$$

where

$$b = \max_{i_N} \left\{ \sum_{i_2 \sim i_{N-1}} \max_{i_1} [c(u_{i_1}^1, u_{i_2}^2)] \prod_{m=2}^{N-1} c(u_{i_m}^m, u_{i_{m+1}}^{m+1}) \right\}. \quad (4)$$

and $S(\rho)$ is the *von Neumann entropy* of the quantum state ρ . Thus the state-independent uncertainty relation for multi-measurement is the corresponding inequality by ignoring $S(\rho)$. In fact, the state-independent inequality generalizes Maassen-Uffink's bound, which suggests that there are rooms for improvement in regards to Coles-Piani's bound. Such an improvement will be useful for further applications in quantum information processing, especially in quantum cryptography when several measurements are present. For the importance of entropic

uncertainty relations and other applications, the reader is referred to [8, 9].

The aim of this article is to find several tighter bounds for multi-measurements in comparison with the bound of Eq.(3) by using majorization theory and symmetry. Of course it is a combinatorial or mathematical exercise to obtain bounds for multi-measurements based on the usual entropic sum of two measurements. However, what we will show is that deeper analysis is needed for nontrivial and tighter bounds for multi-measurements, and applications of majorization theory and symmetry inside the physical construction help to obtain true generalization for multi-measurements.

Indeed, from the construction of the universal uncertainty relation [10, 11], the joint probability distribution in vector $P^1 \otimes P^2$, with respect to the measurement M_1 and M_2 , should be controlled by a bound ω that quantifies its uncertainty in terms of majorization and is also independent of the state ρ . Thus, $H(P^1) + H(P^2) \geq H(\omega)$ for any nonnegative Schur concave function H such as the Shannon entropy. Therefore, the generalized universal uncertainty relation for N measurements

$$\bigotimes_{m=1}^N P^m \prec \omega$$

can imply that $\sum_{m=1}^N H(P^m) \geq H(\omega)$ for multi-measurements. In section II, we first give a precise formula of majorization bound for N probability distributions, and discuss two simple forms of the majorization bounds for multi-measurements in connection with Eq. (3). Comparison of our bounds with previously ones in figure 1 shows that our bounds are tighter.

Further study shows that the simple sum of the uncertainties does not completely reveal the physical meaning of the entropic bounds. The reason is that when one computes the sum of the entropies such as Eq. (3), the mathematical summation does not really provide physically correct answer, as the measurement outcomes clearly do not know which order we perform the measurements, and the bound for N -measurement should be independent from the order of measuring. Therefore one should consider the average of all possible orders of measurements. But this average is cumbersome and does not

provide good enough result.

In order to solve this and get operational formulas for the entropic uncertainty relation of multi-measurements, we study the effects of symmetry on majorization bounds in section III and find that there is a large invariant subgroup of the full symmetry group under the action on certain products of probability distribution vectors and logarithms of remaining distributions. After factoring out this invariant factor we obtain a simple average to give our main result in Section III:

$$\sum_{m=1}^N H(M_m) + (1 - N)S(\rho) \geq -\frac{1}{N}\omega\mathfrak{B}. \quad (5)$$

where ω is the universal majorization bound of N -measurements and \mathfrak{B} is certain vector of logarithmic distributions (cf. Theorem 3). We call this bound an *admixture bound*, since it is obtained by mixing the universal bound from tensor products and factoring out the action of the invariant subgroup of the symmetric group. We then show that this admixture bound is tighter than all previously known bounds in the last part of the section. The exact comparison is charted in figure 2.

II. UNIVERSAL BOUNDS OF MAJORIZATION

Majorization characterizes a balanced partial relationship between two vectors that are comparable and was studied long ago in algebra and analysis. It has been used to study entropic uncertainty relations [12, 13] and played an important role in formulation of state-independent entropic uncertainty relations [10, 11, 14]. A vector x is majorized by another vector y in \mathbb{R}^d : $x \prec y$ if $\sum_{i=1}^k x_i^\downarrow \leq \sum_{i=1}^k y_i^\downarrow$ ($k = 1, 2, \dots, d-1$) and $\sum_{i=1}^d x_i^\downarrow = \sum_{i=1}^d y_i^\downarrow$, where the down-arrow denotes that the components are ordered in decreasing order $x_1^\downarrow \geq \dots \geq x_d^\downarrow$. A nonnegative Schur concave function Φ on \mathbb{R}^d preserves the partial order in the sense that $x \prec y$ implies that $\Phi(x) \geq \Phi(y)$. We adopt the convention to write a probability distribution vector in a short form by omitting the string of zeroes at the end, for example, $(0.6, 0.4, 0, \dots, 0) = (0.6, 0.4)$ and the actual dimension of the vector should be clear from the context.

The tensor product $x \otimes y$ of two vectors $x =$

(x_1, \dots, x_{d_1}) and $y = (y_1, \dots, y_{d_2})$ is defined as $(x_1 y_1, \dots, x_1 y_{d_2}, \dots, x_{d_1} y_1, \dots, x_{d_1} y_{d_2})$, and multi-tensors are defined by associativity. It is well-known that Shannon, Rényi and Tsallis entropies are nonnegative Schur-concave, thus for probability distributions P^1 and P^2 with $P^1 \otimes P^2 \prec \omega$ implies that $\Phi(P^1 \otimes P^2) \geq \Phi(\omega)$ for any of the entropies Φ .

A majorization uncertainty relation for two measurements was well studied in [10, 11]. We now construct the analogous universal upper bound for multi-measurements. Let ρ be a mixed quantum state on a d -dimensional Hilbert space $\mathcal{H} \cong \mathbb{C}^d$, and let M_m ($m = 1, 2, \dots, N$) be N measurements. Assume that M_m has a set of orthonormal eigenvectors $\{|u_{i_m}^m\rangle\}$ ($i_m = 1, 2, \dots, d$), and denote by $P^m = (p_{i_m}^m)$, where $p_{i_m}^m = \langle u_{i_m}^m | \rho | u_{i_m}^m \rangle$ the probability distributions obtained by measuring ρ with respect to bases $\{|u_{i_m}^m\rangle\}$. We can derive a state-independent bound of $\bigotimes_m P^m$ under majorization

$$\bigotimes_{m=1}^N P^m \prec \omega, \quad (6)$$

where the quantity on the left-hand side represents the joint probability distribution induced by measuring ρ with measurements M_m ($m = 1, 2, \dots, N$).

For subsets $\{|u_{i_1}^1\rangle, \dots, |u_{i_{S_1}}^1\rangle\}$, $\{|u_{j_1}^2\rangle, \dots, |u_{j_{S_2}}^2\rangle\}$, \dots , $\{|u_{i_1}^N\rangle, \dots, |u_{i_{S_N}}^N\rangle\}$ of the orthonormal bases of M^1, M^2, \dots, M^N respectively such that $S_1 + S_2 + \dots + S_N = k + N - 1$, we define the matrices $U_{ij}(S_i, S_j)$

$$U_{12}(S_1, S_2) = \begin{pmatrix} \langle u_{i_1}^1 | \\ \langle u_{i_2}^1 | \\ \vdots \\ \langle u_{i_{S_1}}^1 | \end{pmatrix} \cdot \left(|u_{j_1}^2\rangle, |u_{j_2}^2\rangle, \dots, |u_{j_{S_2}}^2\rangle \right) \\ = \begin{pmatrix} \langle u_{i_1}^1 | u_{j_1}^2 \rangle & \langle u_{i_1}^1 | u_{j_2}^2 \rangle & \cdots & \langle u_{i_1}^1 | u_{j_{S_2}}^2 \rangle \\ \langle u_{i_2}^1 | u_{j_1}^2 \rangle & \langle u_{i_2}^1 | u_{j_2}^2 \rangle & \cdots & \langle u_{i_2}^1 | u_{j_{S_2}}^2 \rangle \\ \vdots & \vdots & \ddots & \vdots \\ \langle u_{i_{S_1}}^1 | u_{j_1}^2 \rangle & \langle u_{i_{S_1}}^1 | u_{j_2}^2 \rangle & \cdots & \langle u_{i_{S_1}}^1 | u_{j_{S_2}}^2 \rangle \end{pmatrix}. \quad (7)$$

For simplicity we abbreviate $U_{12}(S_1, S_2)$ by U_{12} . Then $U_{13}, U_{14}, \dots, U_{N-1, N}$ are constructed similarly. We define

the block matrix

$$U(S_1, S_2, \dots, S_N) = \begin{pmatrix} I_{S_1} & U_{12} & \cdots & U_{1N} \\ U_{21} & I_{S_2} & \cdots & U_{2N} \\ \vdots & \vdots & \ddots & \vdots \\ U_{N1} & U_{N2} & \cdots & I_{S_N} \end{pmatrix}. \quad (8)$$

Since the eigenvalues of a Hermitian matrix are real, we adopt the convention to label the eigenvalues in decreasing order. Let $\lambda_1(\bullet)$ and $\sigma_1(\bullet)$ denote the maximal eigenvalue and singular value of a matrix respectively. Generalizing the idea of [11, 14], we introduce the elements s_k by

$$s_k = \max_{\sum_{x=1}^N S_x = k + N - 1} \{\lambda_1(U(S_1, S_2, \dots, S_N))\}. \quad (9)$$

We remark that when $N = 2$, Eq. (9) will degenerate to the s_k defined in [11]. Write

$$\Omega_k = \left(\frac{s_k}{N}\right)^N, \quad (10)$$

then we have $\Omega_1 \leq \Omega_2 \leq \dots \leq \Omega_a < 1$ for some integer $a \leq d^N - 1$ with $\Omega_{a+1} = 1$. With this preparation we can state our universal upper bound for multi-measurements:

Theorem 1. *For any d -dimensional quantum state ρ and N measurements M_m with their probability distributions P^m , we have*

$$\bigotimes_{m=1}^N P^m \prec \omega, \quad (11)$$

where

$$\omega = (\Omega_1, \Omega_2 - \Omega_1, \dots, 1 - \Omega_a). \quad (12)$$

with a being the smallest index such that $\Omega_{a+1} = 1$. Here we have used the short form of the d^N -dimensional vector ω .

Theorem 1 is a generalization of the majorization bound for a pair of two measurements [10, 11]. Due to its key role in our discussion, we include a detailed proof.

Proof of Theorem 1. Consider sums of k elements from the vector $\bigotimes_{m=1}^N P^m$, then they are bounded as follows.

$$\begin{aligned} & (p_{i_1}^1 p_{j_1}^2 \cdots p_{l_1}^N) + \cdots + (p_{i_k}^1 p_{j_k}^2 \cdots p_{l_k}^N) \\ & \leq \max_{S_1 + \cdots + S_N = k + N - 1} \left(\sum_{x=1}^{S_1} \tilde{p}_x^1 \right) \left(\sum_{x=1}^{S_2} \tilde{p}_x^2 \right) \cdots \left(\sum_{x=1}^{S_N} \tilde{p}_x^N \right), \quad (13) \end{aligned}$$

where $\tilde{p}_1^i, \tilde{p}_2^i, \dots, \tilde{p}_{S_i}^i$ are the greatest S_i elements of p_x^i .

Since the arithmetic mean is at least as large as the geometric mean, we derive that

$$\left(\sum_{x=1}^{S_1} \tilde{p}_x^1\right) \left(\sum_{x=1}^{S_2} \tilde{p}_x^2\right) \cdots \left(\sum_{x=1}^{S_N} \tilde{p}_x^N\right) \leq \left(\frac{\sum_{x=1}^{S_1} \tilde{p}_x^1 + \cdots + \sum_{x=1}^{S_N} \tilde{p}_x^N}{N}\right)^N, \quad (14)$$

On the other hand,

$$\begin{aligned} & \sum_{x=1}^{S_1} \tilde{p}_x^1 + \cdots + \sum_{x=1}^{S_N} \tilde{p}_x^N \\ & \leq \max_{\sum_{x=1}^N S_x = k+N-1} \{\lambda_1(U(S_1, S_2, \dots, S_N))\} = s_k, \end{aligned} \quad (15)$$

so we finally get the following estimate:

$$\bigotimes_{m=1}^N P^m \prec (\Omega_1, \Omega_2 - \Omega_1, \dots, 1 - \Omega_a), \quad (16)$$

where $\Omega_k = \binom{s_k}{N}^N$ and Ω_{a+1} is the first component equal to 1, and this gives the desired majorization bound for multi-measurements.

In the case of higher dimensional quantum state ρ , $\lambda_1(U(S_1, S_2, \dots, S_N))$ becomes hard to calculate. However, one can approximate $\lambda_1(U(S_1, S_2, \dots, S_N))$ by the numerical calculation

$$\lambda_1(U(S_1, S_2, \dots, S_N)) = \max_{|u\rangle} \langle u|U(S_1, S_2, \dots, S_N)|u\rangle, \quad (17)$$

where the maximum runs over unit vectors $|u\rangle$, then the right-hand side of Eq. (17) is a deformation of the well-known *Rayleigh-Ritz ratio*. As the unit ball formed by the vectors is compact, *Weierstraß Theorem* ensures the existence of λ_1 . Here we will give two simple estimates of the majorization bound for multi-measurements. To give the first simple estimation, define $CU(1, 2)$ as

$$CU(S_1, S_2) = \begin{pmatrix} 0 & U_{12} & \cdots & 0 \\ U_{21} & 0 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & 0 \end{pmatrix}. \quad (18)$$

Similarly, we can define $CU(S_i, S_j)$ for any pair of i, j such that $1 \leq i, j \leq d$. Then

$$\begin{aligned} U(S_1, S_2, \dots, S_N) &= I_{N+k-1} + CU(S_1, S_2) + \cdots \\ &+ CU(S_{N-1}, S_N). \end{aligned} \quad (19)$$

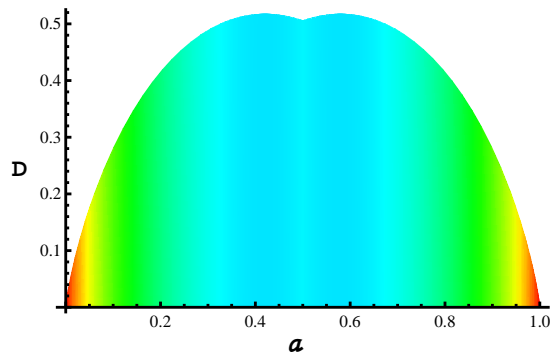


FIG. 1: Difference (D) of $\log \frac{1}{b}$ from $H(\Omega_1, 1 - \Omega_1)$ for $\phi = \pi/2$ with respect to a . The upper curve shows the value of $H(\Omega_1, 1 - \Omega_1) + \log b$ and it is always nonnegative over $0 \leq a \leq 1$.

Using *Weyl's Theorem* on eigenvalues of hermitian matrices, we get that

$$\begin{aligned} & \lambda_1(U(S_1, S_2, \dots, S_N)) \\ &= 1 + \lambda_1(CU(S_1, S_2) + \cdots + CU(S_{N-1}, S_N)) \\ &\leq 1 + \lambda_1(CU(S_1, S_2)) + \cdots + \lambda_1(CU(S_{N-1}, S_N)) \\ &= 1 + \sigma_1(U_{12}) + \cdots + \sigma_1(U_{N-1, N}), \end{aligned} \quad (20)$$

then we define $\hat{\Omega}_k$ by

$$\hat{\Omega}_k = \left(\frac{1 + \hat{s}_k}{N}\right)^N, \quad (21)$$

where

$$\hat{s}_k = \max_{\sum_{x=1}^N S_x = k+N-1} \{\sigma_1(U_{12}) + \cdots + \sigma_1(U_{N-1, N})\}. \quad (22)$$

Therefore we arrive at the following result, which was essentially known in [10].

Theorem 2. For any d -dimensional quantum state ρ and the probability distributions P^m associated to N measurements M_m , we have that

$$\bigotimes_{m=1}^N P^m \prec \hat{\omega}, \quad (23)$$

It is obvious from the construction of $\hat{\omega}$ that the bound is weaker than that of Theorem 1: $\omega \prec \hat{\omega}$.

As for the second approximation, note that the universal bound $\omega \prec (\Omega_1, 1 - \Omega_1)$, which therefore serves as a simple approximation of ω for general N probability distributions. Yet even the bound given by $H(\omega_0)$ with $\omega_0 =$

$(\Omega_1, 1 - \Omega_1)$ outperforms $-\log b$ appeared in Eq. (3). For example, consider three measurements M_i ($i = 1, 2, 3$) in a three-dimensional Hilbert space with eigenvectors $u_1^1 = (1, 0, 0)$, $u_2^1 = (0, 1, 0)$, $u_3^1 = (0, 0, 1)$; $u_1^2 = (\frac{1}{\sqrt{2}}, 0, -\frac{1}{\sqrt{2}})$, $u_2^2 = (0, 1, 0)$, $u_3^2 = (\frac{1}{\sqrt{2}}, 0, \frac{1}{\sqrt{2}})$; $u_1^3 = (\sqrt{a}, e^{i\phi}\sqrt{1-a}, 0)$, $u_2^3 = (\sqrt{1-a}, -e^{i\phi}\sqrt{a}, 0)$ and $u_3^3 = (0, 0, 1)$. With the choice of $\phi = \pi/2$, we see that the simplest majorization bound $\omega_0 = (\Omega_1, 1 - \Omega_1)$ under the Shannon entropy is superior to $-\log b$ over the whole range $0 \leq a \leq 1$, where $b = \max_{i_3} \{\sum_{i_2} \max_{i_1} [c(u_{i_1}^1, u_{i_2}^2)] c(u_{i_2}^2, u_{i_3}^3)\}$. The difference between our second estimation bound $H(\omega_0)$ and Eq. (3), namely $H(\Omega_1, 1 - \Omega_1) + \log b$, is shown in FIG. 1.

III. ADMIXTURE BOUNDS VIA SYMMETRY

As we discussed in the introduction, using Coles and Piani's method, S. Liu *et al.* have given an entropic uncertainty bound for multi-measurements by quantum channels [7]:

$$\sum_{m=1}^N H(M_m) \geq -\log b + (N-1)S(\rho), \quad (24)$$

where

$$b = \max_{i_N} \left\{ \sum_{i_2 \sim i_{N-1}} \max_{i_1} [c(u_{i_1}^1, u_{i_2}^2)] \prod_{m=2}^{N-1} c(u_{i_m}^m, u_{i_{m+1}}^{m+1}) \right\}. \quad (25)$$

We now use the method of symmetry to significantly strengthen the bound. We note that the above bound depends on the order of the measurements, so it is natural to denote the bound as $b(M_1, M_2, \dots, M_N)$ or simply $b(1, 2, \dots, N)$ to specify the order of the measurements M_1, \dots, M_N . Using the apparent symmetry of the measurements, we can define the action of the symmetric group on the bounds. For each permutation $\alpha \in \mathfrak{S}_N$ we define

$$\alpha b(1, \dots, N) = b(\alpha(1), \alpha(2), \dots, \alpha(N)). \quad (26)$$

and observe that \mathfrak{S}_N leaves the second term $(N-1)S(\rho)$ of Eq. (24) invariant. This immediately implies the following entropic uncertainty relation:

$$\sum_{m=1}^N H(M_m) + (1-N)S(\rho) \geq -\log b_{min}, \quad (27)$$

where

$$b_{min} = \min_{\alpha \in \mathfrak{S}_N} \{b(\alpha(1), \alpha(2), \dots, \alpha(N))\}. \quad (28)$$

Apparently $-\log b_{min} \geq -\log b$, so this new bound $-\log b_{min} + (N-1)S(\rho)$ is tighter than the bound appeared in [7]. This shows that the action of the symmetry group can significantly improve the bound. We remark that a similar consideration has been discussed in [15]. Our treatment has clarified how the symmetric group acts on the measurements, which plays an important role in our further investigation.

Now we discuss how to blend the \mathfrak{S}_N -symmetry and the method of quantum channels to derive a tighter bound than we did in the above.

Suppose we are given N measurements M_1, \dots, M_N with orthonormal bases $\{|u_{i_j}^j\rangle\}$. For a multi-index (i_1, \dots, i_N) , where $1 \leq i_j \leq d$, we define the multi-overlap

$$c_{i_1, \dots, i_N}^{1, \dots, N} = c(u_{i_1}^1, u_{i_2}^2) c(u_{i_2}^2, u_{i_3}^3) \cdots c(u_{i_{N-1}}^{N-1}, u_{i_N}^N).$$

Then we have that (cf. [7])

$$\begin{aligned} & (1-N)S(\rho) + \sum_{m=1}^N H(M_m) \\ & \geq -\text{Tr}(\rho \log \sum_{i_1, i_2, \dots, i_N} p_{i_1}^1 c_{i_1, \dots, i_N}^{1, \dots, N} [u_{i_N}^N]) \\ & = -\sum_{i_N} p_{i_N}^N \log \sum_{i_1, i_2, \dots, i_{N-1}} p_{i_1}^1 c_{i_1, \dots, i_{N-1}}^{1, \dots, N} \\ & := I(1, 2, \dots, N), \end{aligned} \quad (29)$$

where $[u]$ stands for $|u\rangle\langle u|$. Note that the above inequality is obtained by a fixing order of M_1, \dots, M_N which explains why we can denote the last expression as $I(1, 2, \dots, N)$. Therefore for any permutation $\alpha \in \mathfrak{S}_N$, one has that

$$\begin{aligned} & (1-N)S(\rho) + \sum_{m=1}^N H(M_m) \\ & \geq I(\alpha(1), \alpha(2), \dots, \alpha(N)), \end{aligned} \quad (30)$$

Taking the average of all permutations, we arrive at the

following relation

$$(1-N)S(\rho) + \sum_{m=1}^N H(M_m) \geq \frac{\sum_{\alpha \in \mathfrak{S}_N} I(\alpha(1), \dots, \alpha(N))}{N!}. \quad (31)$$

Further analysis of the action of the symmetric group on the bound $I(\alpha(1), \dots, \alpha(N))$ shows that only the first and the last indices matter in the formula, as the bound is invariant under the action of any permutation from $\mathfrak{S}_{2, \dots, N-1}$. Among the remaining $N(N-1)$ permutations, it is enough to consider the cyclic group of N permutations. Therefore the above average can be simplified to the following form:

$$(1-N)S(\rho) + \sum_{m=1}^N H(M_m) \geq \frac{\sum_{\text{cyclic } \alpha} I(\alpha(1), \dots, \alpha(N))}{N}, \quad (32)$$

where the sum runs through all N cyclic permutations $(12 \dots N), (23 \dots 1), \dots, (N1 \dots N-1)$.

Let's consider the case of three measurements M_m in detail. By using Eq. (29), we get that

$$-2S(\rho) + \sum_{m=1}^3 H(M_m) \geq -\alpha \left(\sum_{i_3} p_{i_3}^3 \log \sum_{i_1, i_2} p_{i_1}^1 p_{i_2}^2 c_{i_1 i_2 i_3}^{123} \right) := \alpha(I(1, 2, 3)), \quad (33)$$

for any $\alpha \in \mathfrak{S}_3$, thus

$$\begin{aligned} & -2S(\rho) + \sum_{m=1}^3 H(M_m) \\ & \geq \frac{1}{3} (I(1, 2, 3) + I(2, 3, 1) + I(3, 1, 2)) \\ & = \frac{\sum_{i_1, i_2, i_3} p_{i_1}^1 p_{i_2}^2 p_{i_3}^3 \log \sum_{j_1, j_2, j_3} p_{j_1}^1 p_{j_2}^2 p_{j_3}^3 c_{k_1 j_2 i_3}^{123} c_{k_2 j_3 i_1}^{231} c_{k_3 j_1 i_2}^{312}}{-3} \end{aligned} \quad (34)$$

where the sum inside logarithm runs over $j_1, j_2, j_3, k_1, k_2, k_3$. For multi-index (i_1, i_2, i_3) we define the d^3 -dimensional vector $\mathfrak{A}_{i_1, i_2, i_3}$ given by the elements

$$\sum_{j_1, j_2, j_3} c_{k_1 j_2 i_3}^{123} c_{k_2 j_3 i_1}^{231} c_{k_3 j_1 i_2}^{312}, \quad (35)$$

and sorted in decreasing order with respect to multi-indices (k_1, k_2, k_3) (lexicographic order). Combined with

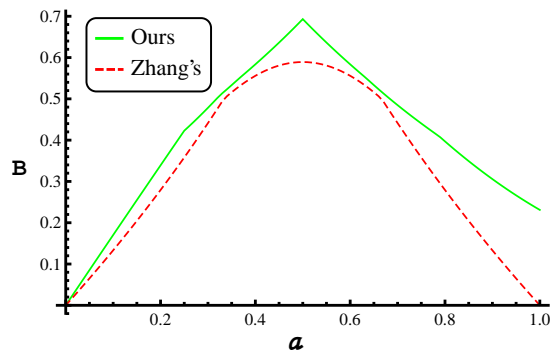


FIG. 2: Comparison of the admixture bound with Zhang *et al.*'s bound for $\phi = \pi/2$ with $a \in [0, 1]$. Our bound in green is shown as the top curve and always tighter. Here \ln is used on the bound axis (B).

the majorization bound $\omega \in \mathbb{R}^{d^3}$ formulated in section II, we immediately get that

$$-\log \sum p_{k_1}^1 p_{k_2}^2 p_{k_3}^3 c_{k_1 j_2 i_3}^{123} c_{k_2 j_3 i_1}^{231} c_{k_3 j_1 i_2}^{312} \geq -\log(\omega \cdot \mathfrak{A}_{i_1, i_2, i_3}). \quad (36)$$

Then we introduce another d^3 -dimensional vector \mathfrak{B} defined by $\mathfrak{B}_{i_1, i_2, i_3} = \log(\omega \cdot \mathfrak{A}_{i_1, i_2, i_3})$ and sorted in decreasing order with respect to multi-indices (i_1, i_2, i_3) in the lexicographic order. Therefore we obtain the following *admixture* bound for 3 measurements

$$-2S(\rho) + \sum_{m=1}^3 H(M_m) \geq -\frac{1}{3} \omega \mathfrak{B}. \quad (37)$$

The new bound provides an improved lower bound for the uncertainty relation. In Fig. 3 we give an example to show that the admixture bound completely outperforms the other bounds that we have known so far for multi-measurements. Moreover, this admixture bound can be easily extended to multi-measurements.

Let $M_i = \{|u_{i_j}^i\rangle\}$ be N measurements, where $i = 1, 2, \dots, N$ and $j = 1, 2, \dots, d$. For each multi-index (i_1, i_2, \dots, i_N) we introduce a d^N -dimensional vector $\mathfrak{A}_{i_1, i_2, \dots, i_N}$ with the entries

$$\sum_{\mathbf{i} \mathbf{j} \dots \mathbf{k}} c_{k_1 j_2 \dots i_N}^{12 \dots N} c_{k_2 j_3 \dots i_1}^{23 \dots 1} \dots c_{k_N j_1 \dots i_{N-1}}^{N1 \dots N-1}$$

where the sum runs over all indices except $\mathbf{i} = (i_1 \dots i_N)$ and $\mathbf{k} = (k_1 \dots k_N)$, and then sorted in decreasing order with respect to lexicographic order of multi-indices

(k_1, \dots, k_N) . Set $\log(\omega \cdot \mathfrak{A}_{i_1, i_2, \dots, i_N}) := \mathfrak{B}_{i_1, i_2, \dots, i_N}$ as the next d^N -dimensional vector with ω being the majorization bound for N measurements formulated in the section II. Here $\mathfrak{B}_{i_1, i_2, \dots, i_N}$ is assumed to be arranged in decreasing order with respect to the multi-indices (i_1, i_2, \dots, i_N) lexicographically. The following result is then proved similarly as before.

Theorem 3. *The following entropic uncertainty relation holds,*

$$\sum_{m=1}^N H(M_m) + (1 - N)S(\rho) \geq -\frac{1}{N}\omega\mathfrak{B}. \quad (38)$$

The admixture bound is tighter than the previously known bounds. In fact, Fig. 2 depicts a comparison of our bound with that of J. Zhang *et al.* [15], while the latter is known to be tighter than the bound appeared in [7].

IV. DISCUSSION

In this paper, we have derived several tighter bounds for entropic uncertainty relations of multi-measurements and in particular an admixture bound is obtained and proved to be tighter than all previously known bounds. Inspired by the recent work [6, 7, 10, 11, 14] we have taken the advantage of unitary matrix $U(S_1, S_2, \dots, S_N)$ and come up with the universal bound for the multi-tensor products of distribution vectors. To derive a deeper and better bound for N measurements, we have studied the action of the symmetric group \mathfrak{S}_N in combination with the universal vector bound of the distribution vectors and quantum channels. The derived admixture bound turns out to be non-trivial bound for the uncertainties of N measurements. Detailed comparisons with previously known bounds are given in figures, and our admixture bound seems to outperform the other bounds most of the time.

Entropy characterizes and quantifies the physical essence of information resources in a mathematical manner. The computational and operational properties of entropy make entropic uncertainty relations useful for quantum key distributions and other quantum cryptography tasks, which can be performed relatively easy in

a physical laboratory. Our new bounds are expected to be useful in handling large data for these and further quantum information processings.

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