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systems

by

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Polygamy relations of multipartite systems

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We investigate the polygamy relations of multipartite quantum states. General polygamy inequalities are given in the α th ($\alpha \geq 2$) power of concurrence of assistance, β th ($\beta \geq 1$) power of entanglement of assistance, and the squared convex-roof extended negativity of assistance (SCRENoA).

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INTRODUCTION

Quantum entanglement is an important kind of quantum correlation, plays essential roles in quantum information processing [1–8]. One of the fundamental differences between classical and quantum correlations lies on the sharability among the subsystems. Different from the classical correlation, quantum correlation cannot be freely shared. Monogamy relation is important in the sense that it gives rise to the distribution of correlation in the multipartite quantum system and has a unique feature of keeping security in quantum key distribution [9].

For the systems of three qubits, a kind of monogamy of bipartite quantum entanglement in concurrence [10] can be described by Coffman-Kundu-Wootters CKW inequality [11], $\mathcal{E}_{A|BC} \geq \mathcal{E}_{AB} + \mathcal{E}_{AC}$, where $\mathcal{E}_{A|BC}$ denotes the entanglement between systems A and BC . Whereas monogamy of entanglement shows the restricted sharability of multipartite entanglement, the distribution of entanglement, or entanglement of assistance [12], in multipartite quantum systems was shown to have a dually monogamous (polygamous) property. Note that the monogamy of entanglement inequalities provide an upper bound for bipartite sharability of entanglement in a multipartite system, and the same quantity sets a lower bound for the distribution of bipartite entanglement in a multipartite system, i.e., $E_{aA|BC} \leq E_{aAB} + E_{aAC}$ for a tripartite quantum state ρ_{ABC} , where $E_{aA|BC}$ is the assisted entanglement [12] between A and BC . The polygamy inequality was first obtained in terms of the tangle of assistance [12] among three-qubit systems, and it was generalized to the multiqubit system with the help of additional entanglement measures [13–15]. In [16–18], people derived a general polygamy inequality of multipartite entanglement beyond qubit based on the entanglement of assistance.

Recently, monogamy and polygamy relations of multiqubit entanglement have been studied in terms of non-negative power of entanglement measures and assisted

entanglement measures. In [19–21], the authors have shown that the x th power of the entanglement of formation ($x \geq \sqrt{2}$) and the concurrence ($x \geq 2$) satisfy multiqubit monogamy inequalities. Monogamy relations for quantum steering have also been demonstrated in [22–26]. Later, polygamy inequalities were also proposed in terms of α th ($0 \leq \alpha \leq 1$) power of square of convex-roof extended negativity (SCREN) and the entanglement of assistance [27, 28]. In [29], the authors introduced a definition of polygamy relations without inequalities. However, it is still not clear for the polygamy relation of the concurrence of assistance τ_a^α ($\alpha \geq 2$) and the β th ($\beta \geq 1$) power of entanglement of assistance E_a^β and the SCREN of assistance (SCRENoA) $(N_{sc}^a)^\beta$. In this paper, we study the general polygamy inequalities of τ_a^α , E_a^β and $(N_{sc}^a)^\beta$ with $\alpha \geq 2$ and $\beta \geq 1$, respectively.

We first recall monogamy and polygamy inequalities related to concurrence and concurrence of assistance. Let \mathbb{H}_X denote a discrete finite-dimensional complex vector space associated with a quantum subsystem X . For a bipartite pure state $|\psi\rangle_{AB} \in \mathbb{H}_A \otimes \mathbb{H}_B$, the concurrence is given by [30–32], $C(|\psi\rangle_{AB}) = \sqrt{2[1 - \text{Tr}(\rho_A^2)]}$, where ρ_A is the reduced density matrix obtained by tracing over the subsystem B , $\rho_A = \text{Tr}_B(|\psi\rangle_{AB}\langle\psi|)$. The concurrence for a bipartite mixed state ρ_{AB} is defined by the convex roof extension, $C(\rho_{AB}) = \min_{\{p_i, |\psi_i\rangle\}} \sum_i p_i C(|\psi_i\rangle)$, where the minimum is taken over all possible pure state decompositions of $\rho_{AB} = \sum_i p_i |\psi_i\rangle\langle\psi_i|$, with $p_i \geq 0$, $\sum_i p_i = 1$ and $|\psi_i\rangle \in \mathbb{H}_A \otimes \mathbb{H}_B$.

For a tripartite state $|\psi\rangle_{ABC}$, the concurrence of assistance is defined by [33, 34], $C_a(|\psi\rangle_{ABC}) \equiv C_a(\rho_{AB}) = \max_{\{p_i, |\psi_i\rangle\}} \sum_i p_i C(|\psi_i\rangle)$, where the maximum is taken over all possible pure state decompositions of $\rho_{AB} = \text{Tr}_C(|\psi\rangle_{ABC}\langle\psi|) = \sum_i p_i |\psi_i\rangle_{AB}\langle\psi_i|$. For pure states $\rho_{AB} = |\psi\rangle_{AB}\langle\psi|$, one has $C(|\psi\rangle_{AB}) = C_a(\rho_{AB})$.

For an N -qubit state $\rho_{AB_1 \dots B_{N-1}} \in \mathbb{H}_A \otimes \mathbb{H}_{B_1} \otimes \dots \otimes \mathbb{H}_{B_{N-1}}$, the concurrence $C(\rho_{A|B_1 \dots B_{N-1}})$ of the state $\rho_{AB_1 \dots B_{N-1}}$, viewed as a bipartite state under the par-

tion A and B_1, B_2, \dots, B_{N-1} , satisfies the Coffman-Kundu-Wootters inequality [35, 36],

$$C^2(\rho_{A|B_1, B_2, \dots, B_{N-1}}) \geq \sum_{i=1}^{N-1} C^2(\rho_{AB_i}), \quad (1)$$

where $\rho_{AB_i} = \text{Tr}_{B_1 \dots B_{i-1} B_{i+1} \dots B_{N-1}}(\rho_{AB_1 \dots B_{N-1}})$. Further improved monogamy relations are presented in [19, 21]. The dual inequality in terms of the concurrence of assistance for N -qubit states has the form [37],

$$C^2(\rho_{A|B_1, B_2, \dots, B_{N-1}}) \leq \sum_{i=1}^{N-1} C_a^2(\rho_{AB_i}). \quad (2)$$

Now, let us consider a bipartite pure state of arbitrary dimension $d_1 \times d_2$, $|\phi\rangle_{AB} = \sum_{i=1}^{d_1} \sum_{k=1}^{d_2} a_{ik} |ik\rangle_{AB}$ in $C^{d_1} \otimes C^{d_2}$. The squared concurrence of $|\phi\rangle_{AB}$ can be expressed as [38]

$$C^2(|\phi\rangle_{AB}) = 2(1 - \text{Tr}(\rho_A^2)) = 4 \sum_{i < j} \sum_{k < l} |a_{ik} a_{jl} - a_{il} a_{jk}|^2. \quad (3)$$

For a mixed state $\rho_{AB} = \sum_i p_i |\phi_i\rangle_{AB} \langle \phi_i|$, its concurrence of assistance satisfies [39]

$$\begin{aligned} C_a(\rho_{AB}) &= \max_{\{p_i, |\phi_i\rangle\}} \sum_i p_i C(|\phi_i\rangle) \\ &\leq \sum_{m=1}^{D_1} \sum_{n=1}^{D_2} (\max_i \sum_i p_i |\langle \phi_i | (L_A^m \otimes L_B^n) | \phi_i^* \rangle|) \\ &= \sum_{m=1}^{D_1} \sum_{n=1}^{D_2} C_a((\rho_{AB})_{mn}) := \tau_a(\rho_{AB}), \end{aligned} \quad (4)$$

where

$$D_1 = d_1(d_1 - 1)/2, \quad D_2 = d_2(d_2 - 1)/2, \quad (5)$$

$$L_A^m = P_A^m(-|i\rangle_A \langle j| + |j\rangle_A \langle i|) P_A^m, \quad (6)$$

$$L_B^n = P_B^n(-|k\rangle_B \langle l| + |l\rangle_B \langle k|) P_B^n \quad (7)$$

with $P_A^m = |i\rangle_A \langle i| + |j\rangle_A \langle j|$ and $P_B^n = |k\rangle_B \langle k| + |l\rangle_B \langle l|$ being the projectors to the subspaces spanned by $\{|i\rangle_A, |j\rangle_A\}$ and $\{|k\rangle_B, |l\rangle_B\}$, respectively. A general polygamy inequality for any multipartite pure state $|\phi\rangle_{A_1 \dots A_n} \in C^{d_1} \otimes \dots \otimes C^{d_n}$ was established as [39],

$$\tau_a^2(|\phi\rangle_{A_1 A_2 \dots A_n}) \leq \sum_{i=2}^n \tau_a^2(\rho_{A_1 A_i}), \quad (8)$$

where $\rho_{A_1 A_k}$ is the reduced density matrix $|\phi\rangle_{A_1 A_2 \dots A_n}$ with respect to subsystem $A_1 A_k$, $k = 2, \dots, n$.

POLYGAMY RELATION FOR CONCURRENCE OF ASSISTANCE

[Lemma 1]. For any real numbers x and t , $t \geq 1$, $x \geq 1$, we have $(1+t)^x \leq 1 + (2^x - 1)t^x$.

[Proof]. Let $f(x, y) = (1+y)^x - y^x$ with $x \geq 1$, $0 < y \leq 1$, $\frac{\partial f}{\partial y} = x[(1+y)^{x-1} - y^{x-1}] \geq 0$. Therefore, $f(x, y)$ is an increasing function of y , i.e., $f(x, y) \leq f(x, 1) = 2^x - 1$. Set $y = \frac{1}{t}$, $t \geq 1$. We obtain $(1+t)^x \leq 1 + (2^x - 1)t^x$. Notice when $t = 1$, the inequality is true. \square

The following theorem provides a class of polygamy inequalities satisfied by the α -power of τ_a . For convenience, we denote $\tau_a(\rho_{AB_i}) = \tau_{aAB_i}$ the concurrence of assistance ρ_{AB_i} and $\tau_a(\rho_{A|B_0 B_1 \dots B_{N-1}}) = \tau_{aA|B_0 B_1 \dots B_{N-1}}$.

[Theorem 1]. For any tripartite pure state $\rho_{ABC} \in H_A \otimes H_B \otimes H_C$:

(1) if $\tau_{aAB} \geq \tau_{aAC}$, the concurrence of assistance satisfies

$$\tau_{aA|BC}^\alpha \leq \tau_{aAC}^\alpha + (2^{\frac{\alpha}{2}} - 1)\tau_{aAB}^\alpha \quad (9)$$

for $\alpha \geq 2$.

(2) if $\tau_{aAB} \leq \tau_{aAC}$, the concurrence of assistance satisfies

$$\tau_{aA|BC}^\alpha \leq \tau_{aAB}^\alpha + (2^{\frac{\alpha}{2}} - 1)\tau_{aAC}^\alpha \quad (10)$$

for $\alpha \geq 2$.

[Proof]. For arbitrary tripartite pure state ρ_{ABC} , one has [39], $\tau_{aA|BC}^2 \leq \tau_{aAC}^2 + \tau_{aAB}^2$. If $\tau_{aAB} (\tau_{aAC}) = 0$, the inequality (9) or (10) are true obviously. Therefore, assuming $\tau_{aAB} \geq \tau_{aAC} > 0$, we have

$$\begin{aligned} \tau_{aA|BC}^{2x} &\leq (\tau_{aAB}^2 + \tau_{aAC}^2)^x \\ &= \tau_{aAC}^{2x} \left(1 + \frac{\tau_{aAB}^2}{\tau_{aAC}^2}\right)^x \\ &\geq \tau_{aAC}^{2x} \left(1 + (2^x - 1) \left(\frac{\tau_{aAB}^2}{\tau_{aAC}^2}\right)^x\right) \\ &= \tau_{aAC}^{2x} + (2^x - 1)\tau_{aAB}^{2x}, \end{aligned} \quad (11)$$

where the second inequality is true due to the inequality $(1+t)^x \leq 1 + (2^x - 1)t^x$ for $x \geq 1$ and $t = \frac{\tau_{aAB}^2}{\tau_{aAC}^2} \geq 1$. Denote $2x = \alpha$. We obtain $\alpha \geq 2$ as $x \geq 1$. Then we have the inequality (9). If $\tau_{aAB} \leq \tau_{aAC}$, Similarly we get (10).

Example 1. Let us consider the three-qubit state $|\psi\rangle$ in the generalized Schmidt decomposition form,

$$\begin{aligned} |\psi\rangle &= \lambda_0 |000\rangle + \lambda_1 e^{i\varphi} |100\rangle + \lambda_2 |101\rangle \\ &\quad + \lambda_3 |110\rangle + \lambda_4 |111\rangle, \end{aligned} \quad (12)$$

where $\lambda_i \geq 0$, $i = 0, 1, 2, 3, 4$ and $\sum_{i=0}^4 \lambda_i^2 = 1$. We have

$\tau_{aA|BC} = 2\lambda_0 \sqrt{\lambda_2^2 + \lambda_3^2 + \lambda_4^2}$, $\tau_{aAB} = 2\lambda_0 \sqrt{\lambda_2^2 + \lambda_4^2}$, and $\tau_{aAC} = 2\lambda_0 \sqrt{\lambda_3^2 + \lambda_4^2}$. Without loss of generality, we set $\lambda_0 = \cos \theta_0$, $\lambda_1 = \sin \theta_0 \cos \theta_1$, $\lambda_2 = \sin \theta_0 \sin \theta_1 \cos \theta_2$, $\lambda_3 = \sin \theta_0 \sin \theta_1 \sin \theta_2 \cos \theta_3$, and $\lambda_4 = \sin \theta_0 \sin \theta_1 \sin \theta_2 \sin \theta_3$, $\theta_i \in [0, \frac{\pi}{2}]$.

For $\lambda_3 \geq \lambda_2$, i.e. $\tau_{aAC} \geq \tau_{aAB}$:

(a) if $\theta_2 = \frac{\pi}{2}$,

$$\begin{aligned} & \tau_{a|ABC}^\alpha - \tau_{aAB}^\alpha - (2^{\frac{\alpha}{2}} - 1)\tau_{aAC}^\alpha \\ &= (2\lambda_0)^\alpha \left[(\lambda_2^2 + \lambda_3^2 + \lambda_4^2)^{\frac{\alpha}{2}} - (\lambda_2^2 + \lambda_4^2)^{\frac{\alpha}{2}} \right. \\ & \quad \left. - (2^{\frac{\alpha}{2}} - 1)(\lambda_3^2 + \lambda_4^2)^{\frac{\alpha}{2}} \right] \\ &= 2^\alpha \cos^\alpha \theta_0 \sin^\alpha \theta_0 \sin^\alpha \theta_1 (2 - \sin^\alpha \theta_3 - 2^{\frac{\alpha}{2}}) \\ &\leq 0, \end{aligned}$$

where $\alpha \geq 2$ and the inequality is due to $\sin \theta_3 \geq 0$.

(b) If $\theta_2 \neq \frac{\pi}{2}$, we denote $t_1 = \frac{\sin^2 \theta_2}{\cos^2 \theta_2} \geq 1$, then we have

$$\begin{aligned} & \tau_{a|ABC}^\alpha - \tau_{aAB}^\alpha - (2^{\frac{\alpha}{2}} - 1)\tau_{aAC}^\alpha \\ &= (2\lambda_0)^\alpha \left[(\lambda_2^2 + \lambda_3^2 + \lambda_4^2)^{\frac{\alpha}{2}} - (\lambda_2^2 + \lambda_4^2)^{\frac{\alpha}{2}} \right. \\ & \quad \left. - (2^{\frac{\alpha}{2}} - 1)(\lambda_3^2 + \lambda_4^2)^{\frac{\alpha}{2}} \right] \\ &= 2^\alpha \cos^\alpha \theta_0 \sin^\alpha \theta_0 \sin^\alpha \theta_1 \left[1 - (\cos^2 \theta_2 + \sin^2 \theta_2 \sin^2 \theta_3)^{\frac{\alpha}{2}} \right. \\ & \quad \left. - (2^{\frac{\alpha}{2}} - 1) \sin^\alpha \theta_2 \right] \\ &\leq 2^\alpha \cos^\alpha \theta_0 \sin^\alpha \theta_0 \sin^\alpha \theta_1 \left[1 - \cos^\alpha \theta_2 - (2^{\frac{\alpha}{2}} - 1) \sin^\alpha \theta_2 \right] \\ &= 2^\alpha \cos^\alpha \theta_0 \sin^\alpha \theta_0 \sin^\alpha \theta_1 \left[1 - \cos^\alpha \theta_2 \left(1 + (2^{\frac{\alpha}{2}} - 1)t_1^{\frac{\alpha}{2}} \right) \right] \\ &\leq 2^\alpha \cos^\alpha \theta_0 \sin^\alpha \theta_0 \sin^\alpha \theta_1 \left[1 - \cos^\alpha \theta_2 (1 + t_1)^{\frac{\alpha}{2}} \right] \\ &= 2^\alpha \cos^\alpha \theta_0 \sin^\alpha \theta_0 \sin^\alpha \theta_1 \left[1 - \cos^\alpha \theta_2 \left(1 + \frac{\sin^2 \theta_2}{\cos^2 \theta_2} \right)^{\frac{\alpha}{2}} \right] \\ &= 0, \end{aligned}$$

where $\alpha \geq 2$ and the second inequality is due to Lemma 1.

Therefore, we have $\tau_{a|ABC}^\alpha \leq \tau_{aAB}^\alpha + (2^{\frac{\alpha}{2}} - 1)\tau_{aAC}^\alpha$ for $\alpha \geq 2$.

When $\lambda_3 \leq \lambda_2$, i.e. $\tau_{aAC} \leq \tau_{aAB}$, from similar analysis we can obtain $\tau_{a|ABC}^\alpha \leq \tau_{aAC}^\alpha + (2^{\frac{\alpha}{2}} - 1)\tau_{aAB}^\alpha$ for $\alpha \geq 2$.

Specially, when $\theta_2 = \frac{\pi}{2}, \alpha = 2, \theta_3 = 0$, i.e. $|\psi\rangle = \cos \theta_0 |000\rangle + \sin \theta_0 \cos \theta_1 e^{i\varphi} |100\rangle + \sin \theta_0 \sin \theta_1 |110\rangle$, the inequality in (9) is saturated. Generalizing the conclusion in Theorem 1 to N partite case, we have the following result.

[Theorem 2]. For any multipartite pure state $\rho_{AB_0 \dots B_{N-1}}$, if $\tau_{aAB_i}^2 \leq \sum_{k=i+1}^{N-1} \tau_{aAB_k}^2$ for $i = 0, 1, \dots, m$, and $\tau_{aAB_j}^2 \geq \sum_{k=j+1}^{N-1} \tau_{aAB_k}^2$ for $j = m + 1, \dots, N - 2, \forall 1 \leq m \leq N - 3, N \geq 4$, we have

$$\begin{aligned} & \tau_{a|B_0 B_1 \dots B_{N-1}}^\alpha \leq \\ & \tau_{aAB_0}^\alpha + (2^{\frac{\alpha}{2}} - 1)\tau_{aAB_1}^\alpha + \dots + (2^{\frac{\alpha}{2}} - 1)^m \tau_{aAB_m}^\alpha \\ & + (2^{\frac{\alpha}{2}} - 1)^{m+2} (\tau_{aAB_{m+1}}^\alpha + \dots + \tau_{aAB_{N-2}}^\alpha) \\ & + (2^{\frac{\alpha}{2}} - 1)^{m+1} \tau_{aAB_{N-1}}^\alpha, \end{aligned} \quad (13)$$

for $\alpha \geq 2$.

[Proof]. From the inequality (8) and Theorem 1, we have

$$\begin{aligned} & \tau_{a|B_0 B_1 \dots B_{N-1}}^\alpha \\ & \leq \tau_{aAB_0}^\alpha + (2^{\frac{\alpha}{2}} - 1) \left(\sum_{i=1}^{N-1} \tau_{aAB_i}^2 \right)^{\frac{\alpha}{2}} \\ & \leq \tau_{aAB_0}^\alpha + (2^{\frac{\alpha}{2}} - 1)\tau_{aAB_1}^\alpha + (2^{\frac{\alpha}{2}} - 1)^2 \left(\sum_{i=2}^{N-1} \tau_{aAB_i}^2 \right)^{\frac{\alpha}{2}} \\ & \leq \dots \\ & \leq \tau_{aAB_0}^\alpha + (2^{\frac{\alpha}{2}} - 1)\tau_{aAB_1}^\alpha + \dots + (2^{\frac{\alpha}{2}} - 1)^m \tau_{aAB_m}^\alpha \\ & \quad + (2^{\frac{\alpha}{2}} - 1)^{m+1} \left(\sum_{i=m+1}^{N-1} \tau_{aAB_i}^2 \right)^{\frac{\alpha}{2}}. \end{aligned} \quad (14)$$

Similarly, as $\tau_{aAB_j}^2 \geq \sum_{k=j+1}^{N-1} \tau_{aAB_k}^2$ for $j = m + 1, \dots, N - 2$, we get

$$\begin{aligned} & \left(\sum_{i=m+1}^{N-1} \tau_{aAB_i}^2 \right)^{\frac{\alpha}{2}} \\ & \leq (2^{\frac{\alpha}{2}} - 1)\tau_{aAB_{m+1}}^\alpha + \left(\sum_{i=m+2}^{N-1} \tau_{aAB_i}^2 \right)^{\frac{\alpha}{2}} \\ & \leq (2^{\frac{\alpha}{2}} - 1)(\tau_{aAB_{m+1}}^\alpha + \dots + \tau_{aAB_{N-2}}^\alpha) \\ & \quad + \tau_{aAB_{N-1}}^\alpha. \end{aligned} \quad (15)$$

Combining (14) and (15), we have Theorem 2. \square

In Theorem 2, if $\tau_{aAB_i}^2 \leq \sum_{j=i+1}^{N-1} \tau_{aAB_j}^2$ for all $i = 0, 1, \dots, N - 2$, then we have the following conclusion:

[Theorem 3]. For any multipartite pure state $\rho_{AB_0 \dots B_{N-1}}$, if $\tau_{aAB_i}^2 \leq \sum_{j=i+1}^{N-1} \tau_{aAB_j}^2$ for all $i = 0, 1, \dots, N - 2$, we have

$$\tau_{a|B_0 B_1 \dots B_{N-1}}^\alpha \leq \sum_{j=0}^{N-1} (2^{\frac{\alpha}{2}} - 1)^j \tau_{aAB_j}^\alpha, \quad (16)$$

for $\alpha \geq 2$.

Example 2. We consider again the pure state (12). Setting $\lambda_0 = \lambda_1 = \frac{1}{2}, \lambda_2 = \lambda_3 = \lambda_4 = \frac{\sqrt{6}}{6}$, one has $\tau_{a|ABC} = \frac{\sqrt{2}}{2}, \tau_{aAB} = \tau_{aAC} = \frac{\sqrt{3}}{3}$. Let $y = \tau_{aAB}^\alpha + (2^{\frac{\alpha}{2}} - 1)\tau_{aAC}^\alpha - \tau_{a|ABC}^\alpha = 2^{\frac{\alpha}{2}} \left(\frac{\sqrt{3}}{3} \right)^\alpha, \alpha \geq 2$, be the residual concurrence of assistance. From our results, one can see that $y \geq 0$ for $\alpha \geq 2$, which is the case that does not given in [28], see Fig. 1.

POLYGAMY RELATIONS FOR ENTANGLEMENT OF ASSISTANCE

For polygamy inequality beyond qubits, it was shown that the von Neumann entropy can be used to establish a polygamy inequality of tripartite quantum systems [16]. For any arbitrary dimensional tripartite

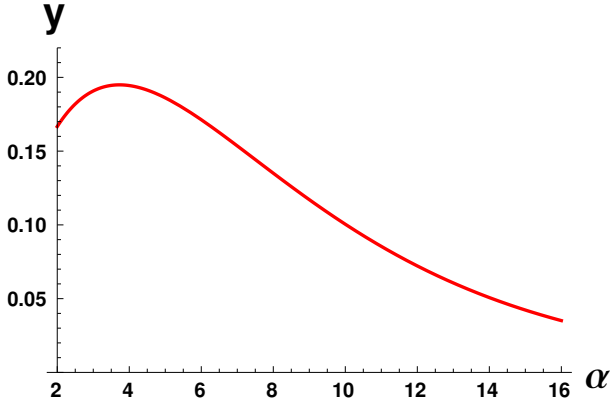


FIG. 1: y is the “residual” entanglement as a function of α with $\alpha \geq 2$.

pure state $|\psi\rangle_{ABC}$, one has $E(|\psi\rangle_{A|BC}) \leq E_a(\rho_{AB}) + E_a(\rho_{AC})$, where $E(|\psi\rangle_{A|BC}) = S(\rho_A)$ is the entropy of entanglement between A and BC in terms of the von Neumann entropy $S(\rho) = -\text{Tr}\rho \ln \rho$, and $E_a(\rho_{AB}) = \max \sum_i p_i E(|\psi_i\rangle_{AB})$, with the maximization taking over all possible pure state decompositions of $\rho_{AB} = \sum_i p_i |\psi_i\rangle_{AB} \langle \psi_i|$. Later, a general polygamy inequality for any multipartite state $\rho_{A_1|A_2 \dots A_n}$ was established, $E_a(\rho_{A_1|A_2 \dots A_n}) \leq \sum_{i=2}^n E_a(\rho_{A_1 A_i})$ [17].

Recently, another class of multipartite polygamy inequalities in terms of the β th power of entanglement of assistance (EOA) has been introduced [28]. For any multipartite state $\rho_{AB_0 B_1 \dots B_{N-1}}$ and $0 \leq \beta \leq 1$, if $E_{aAB_i} \geq \sum_{j=i+1}^{N-1} E_{aAB_j}$ for $i = 0, 1, \dots, N-2$, then $E_{aA|B_0 B_1 \dots B_{N-1}}^\beta \leq \sum_{j=0}^{N-1} \beta^j E_{aAB_j}^\beta$, where $E_a(\rho_{AB_i}) = E_{aAB_i}$ is the entanglement of assistance ρ_{AB_i} and $E_a(\rho_{A|B_0 B_1 \dots B_{N-1}}) = E_{aA|B_0 B_1 \dots B_{N-1}}$. But, for $\beta \geq 1$ the polygamy relations for the β th power of the entanglement of assistance is still not clear.

[Theorem 4]. For any multipartite state $\rho_{AB_0 \dots B_{N-1}}$, if $E_{aAB_i} \leq \sum_{k=i+1}^{N-1} E_{aAB_k}$ for $i = 0, 1, \dots, m$, and $E_{aAB_j} \geq \sum_{k=j+1}^{N-1} E_{aAB_k}$ for $j = m+1, \dots, N-2$, $\forall 1 \leq m \leq N-3$, $N \geq 4$, we have

$$\begin{aligned} E_{aA|B_0 B_1 \dots B_{N-1}}^\beta &\leq \\ E_{aAB_0}^\beta &+ (2^\beta - 1)E_{aAB_1}^\beta + \dots + (2^\beta - 1)^m E_{aAB_m}^\beta \\ &+ (2^\beta - 1)^{m+2} (E_{aAB_{m+1}}^\beta + \dots + E_{aAB_{N-2}}^\beta) \\ &+ (2^\beta - 1)^{m+1} E_{aAB_{N-1}}^\beta, \end{aligned} \quad (17)$$

for $\beta \geq 1$.

[Proof]. From Lemma 1, we have

$$\begin{aligned} E_{aA|B_0 B_1 \dots B_{N-1}}^\beta &\leq E_{aAB_0}^\beta + (2^\beta - 1) \left(\sum_{i=1}^{N-1} E_{aAB_i} \right)^\beta \\ &\leq E_{aAB_0}^\beta + (2^\beta - 1)E_{aAB_1}^\beta + (2^\beta - 1)^2 \left(\sum_{i=2}^{N-1} E_{aAB_i} \right)^\beta \\ &\leq \dots \\ &\leq E_{aAB_0}^\beta + (2^\beta - 1)E_{aAB_1}^\beta + \dots + (2^\beta - 1)^m E_{aAB_m}^\beta \\ &\quad + (2^\beta - 1)^{m+1} \left(\sum_{i=m+1}^{N-1} E_{aAB_i} \right)^\beta. \end{aligned} \quad (18)$$

Similarly, as $E_{aAB_j} \geq \sum_{k=j+1}^{N-1} E_{aAB_k}$ for $j = m+1, \dots, N-2$, we get

$$\begin{aligned} &\left(\sum_{i=m+1}^{N-1} E_{aAB_i} \right)^\beta \\ &\leq (2^\beta - 1)E_{aAB_{m+1}}^\beta + \left(\sum_{i=m+2}^{N-1} E_{aAB_i} \right)^\beta \\ &\leq (2^\beta - 1)(E_{aAB_{m+1}}^\beta + \dots + E_{aAB_{N-2}}^\beta) \\ &\quad + E_{aAB_{N-1}}^\beta. \end{aligned} \quad (19)$$

Combining (18) and (19), we have Theorem 4. \square

As a special case of Theorem 4, if $E_{aAB_i} \leq \sum_{j=i+1}^{N-1} E_{aAB_j}$ for all $i = 0, 1, \dots, N-2$, we have the following conclusion:

[Theorem 5]. For any multipartite state $\rho_{AB_0 \dots B_{N-1}}$, if $E_{aAB_i} \leq \sum_{j=i+1}^{N-1} E_{aAB_j}$ for all $i = 0, 1, \dots, N-2$, we have

$$E_{aA|B_0 B_1 \dots B_{N-1}}^\beta \leq \sum_{j=0}^{N-1} (2^\beta - 1)^j E_{aAB_j}^\beta, \quad (20)$$

for $\beta \geq 1$.

Example 3. Let consider the three-qubit W state $|W\rangle_{ABC} = \frac{1}{\sqrt{3}}(|100\rangle + |010\rangle + |001\rangle)$. We have $E_a(|W\rangle_{A|BC}) = S(\rho_A) = \log_2 3 - \frac{2}{3}$ and $E_a(\rho_{AB}) = E_a(\rho_{AC}) = \frac{2}{3}$. Set $y = E_a^\beta(\rho_{AB}) + (2^\beta - 1)E_a^\beta(\rho_{AC}) - E_a^\beta(|W\rangle_{A|BC}) = 2^\beta (\frac{2}{3})^\beta - (\log_2 3 - \frac{2}{3})^\beta$ to be the residual entanglement of assistance. Fig. 2 shows our polygamy inequality for $\beta \geq 1$.

POLYGAMY RELATIONS FOR SCRENOA

Given a bipartite state ρ_{AB} in $H_A \otimes H_B$, the negativity is defined by [41], $N(\rho_{AB}) = (\|\rho_{AB}^{T_A}\| - 1)/2$, where $\rho_{AB}^{T_A}$ is the partially transposed ρ_{AB} with respect to the subsystem A , $\|X\|$ denotes the trace norm of X , i.e., $\|X\| =$

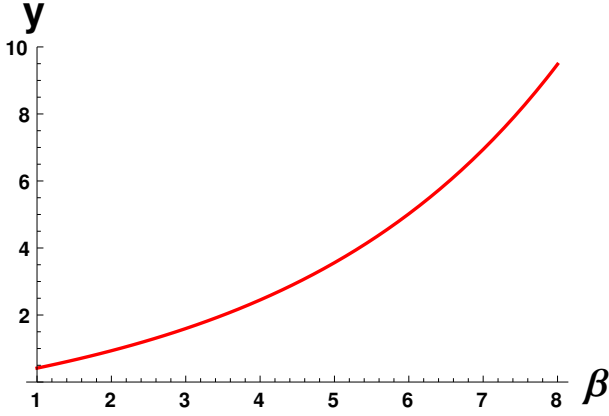


FIG. 2: y is the “residual” entanglement as a function of β with $\beta \geq 1$.

$\text{Tr}\sqrt{XX^\dagger}$. For the purpose of discussion, we use the following definition of negativity, $N(\rho_{AB}) = \|\rho_{AB}^{T_A}\| - 1$. For any bipartite pure state $|\psi\rangle_{AB}$, the negativity $N(\rho_{AB})$ is given by $N(|\psi\rangle_{AB}) = 2 \sum_{i < j} \sqrt{\lambda_i \lambda_j} = (\text{Tr}\sqrt{\rho_A})^2 - 1$, where λ_i are the eigenvalues for the reduced density matrix ρ_A of $|\psi\rangle_{AB}$. For a mixed state ρ_{AB} , the square of convex-roof extended negativity (SCREN) is defined by

$$N_{sc}(\rho_{AB}) = [\min \sum_i p_i N(|\psi_i\rangle_{AB})]^2, \quad (21)$$

where the minimum is taken over all possible pure state decompositions $\{p_i, |\psi_i\rangle_{AB}\}$ of ρ_{AB} . Similar to the duality between concurrence and concurrence of assistance, we also define a dual quantity to SCREN as

$$N_{sc}^a(\rho_{AB}) = [\max \sum_i p_i N(|\psi_i\rangle_{AB})]^2, \quad (22)$$

which we refer to as the SCREN of assistance (SCRENoA), where the maximum is taken over all possible pure state decompositions $\{p_i, |\psi_i\rangle_{AB}\}$ of ρ_{AB} . For convenience, we denote $N_{scAB_i}^a = N_{sc}^a(\rho_{AB_i})$ the SCRENoA of ρ_{AB_i} and $N_{scAB_0 \dots B_{N-1}}^a = N_{sc}^a(|\psi\rangle_{AB_0 \dots B_{N-1}})$.

In [27] it has been shown that $N_{scA|B_0B_1 \dots B_{N-1}}^a \leq \sum_{j=0}^{N-1} N_{scAB_j}^a$. It is further improved that for $0 \leq \beta \leq 1$, $(N_{scA|B_0B_1 \dots B_{N-1}}^a)^\beta \leq \sum_{j=0}^{N-1} \beta^j (N_{scAB_j}^a)^\beta$. But, it is still not clear whether the polygamy relation still holds for the β th ($\beta \geq 1$) power of SCRENoA. With a similar consideration to $\tau_{AB_0 \dots B_{N-1}}$, we have the following result of SCRENoA for $\beta \geq 1$.

[Theorem 6]. For any multipartite state $\rho_{AB_0 \dots B_{N-1}}$, if $N_{scAB_i}^a \leq \sum_{k=i+1}^{N-1} N_{scAB_k}^a$ for $i = 0, 1, \dots, m$, and $N_{scAB_j}^a \geq \sum_{k=j+1}^{N-1} N_{scAB_k}^a$ for $j = m+1, \dots, N-2$, \forall

$1 \leq m \leq N-3$, $N \geq 4$, we have

$$\begin{aligned} (N_{scA|B_0B_1 \dots B_{N-1}}^a)^\beta &\leq (N_{scAB_0}^a)^\beta \\ &+ (2^\beta - 1)(N_{scAB_1}^a)^\beta + \dots + (2^\beta - 1)^m (N_{scAB_m}^a)^\beta \\ &+ (2^\beta - 1)^{m+2} \left((N_{scAB_{m+1}}^a)^\beta + \dots + (N_{scAB_{N-2}}^a)^\beta \right) \\ &+ (2^\beta - 1)^{m+1} (N_{scAB_{N-1}}^a)^\beta, \end{aligned} \quad (23)$$

for $\beta \geq 1$.

[Proof]. From Lemma 1, we have

$$\begin{aligned} E_{aA|B_0B_1 \dots B_{N-1}}^\beta &\leq (N_{scAB_0}^a)^\beta + (2^\beta - 1) \left(\sum_{i=1}^{N-1} N_{scAB_i}^a \right)^\beta \\ &\leq (N_{scAB_0}^a)^\beta + (2^\beta - 1) (N_{scAB_1}^a)^\beta \\ &\quad + (2^\beta - 1)^2 \left(\sum_{i=2}^{N-1} N_{scAB_i}^a \right)^\beta \\ &\leq \dots \\ &\leq (N_{scAB_0}^a)^\beta + (2^\beta - 1) (N_{scAB_1}^a)^\beta + \dots \\ &\quad + (2^\beta - 1)^m (N_{scAB_m}^a)^\beta \\ &\quad + (2^\beta - 1)^{m+1} \left(\sum_{i=m+1}^{N-1} N_{scAB_i}^a \right)^\beta. \end{aligned} \quad (24)$$

Similarly, as $N_{scAB_j}^a \geq \sum_{k=j+1}^{N-1} N_{scAB_k}^a$ for $j = m+1, \dots, N-2$, we get

$$\begin{aligned} \left(\sum_{i=m+1}^{N-1} N_{scAB_i}^a \right)^\beta &\leq (2^\beta - 1) (N_{scAB_{m+1}}^a)^\beta + \left(\sum_{i=m+2}^{N-1} N_{scAB_i}^a \right)^\beta \\ &\leq (2^\beta - 1) \left((N_{scAB_{m+1}}^a)^\beta + \dots + (N_{scAB_{N-1}}^a)^\beta \right) \\ &\quad + (N_{scAB_{N-1}}^a)^\beta. \end{aligned} \quad (25)$$

Combining (24) and (25), we have the Theorem 6. \square

Particularly, the equality in (23) can be established for 4-qubit generalized W -class states $|W\rangle_{AB_1B_2B_3} = \frac{1}{2}(|1000\rangle + |0100\rangle + |0010\rangle + |0001\rangle)$, with $\beta = 1$, which can be seen clearly in example 4 below Theorem 7. Specially, from Theorem 6 we have

[Theorem 7]. For any multipartite state $\rho_{AB_0 \dots B_{N-1}}$, if $E_{aAB_i} \leq \sum_{j=i+1}^{N-1} E_{aAB_j}$ for all $i = 0, 1, \dots, N-2$, we have

$$(N_{scA|B_0B_1 \dots B_{N-1}}^a)^\beta \leq \sum_{j=0}^{N-1} (2^\beta - 1)^j (N_{scAB_j}^a)^\beta \quad (26)$$

for $\beta \geq 1$.

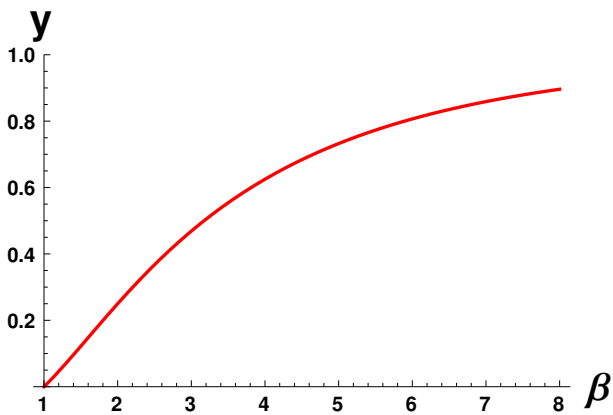


FIG. 3: The “residual” entanglement y as a function of β ($\beta \geq 1$).

Example 4. Let us consider the 4-qubit generalized W -class states,

$$|W\rangle_{AB_1B_2B_3} = \frac{1}{2}(|1000\rangle + |0100\rangle + |0010\rangle + |0001\rangle) \quad (27)$$

We have $N_{scA|B_1B_2B_3}^a = \frac{3}{4}$, $N_{scAB_i}^a = \frac{1}{4}$, $i = 1, 2, 3$. Let y be the difference between the right and left hand of inequality (26). One has $y = [2^\beta + (2^\beta - 1)^2](\frac{1}{4})^\beta - (\frac{3}{4})^\beta$; see Fig. 3.

CONCLUSION

Entanglement monogamy and polygamy are fundamental properties of multipartite entangled states. We have investigated in this work the polygamy relations related to the concurrence of assistance, entanglement of assistance, and SCREN generally for multipartite states. We have found a class of polygamy inequalities of multipartite entanglement in arbitrary-dimensional quantum systems in the α th ($\alpha \geq 2$) power of concurrence of assistance, a case that has not been studied before. Moreover, the β th power of polygamy inequalities have been obtained for the entanglement of assistance and SCREN_{oA} for $\beta \geq 1$. The approach developed in this work is applicable to the study of monogamy properties in other quantum entanglement measures and quantum correlations.

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